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Journal Article

Operator-aware Approach for Boosting Performance in Processing RDF streams

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Agenda



- Overview of RDF Processing
- On Boosting The Processing Throughput
- Challenges in Incremental Evaluation
- Operator-Aware Approach
- Evaluation
- Summary







```
S_{pickup} = \left\{ \begin{array}{l} : ride_1 : taxi : 89...CF4 \\ : ride_1 : pickupTime "2013-01-01 15:11:48". \end{array} \right\}. S_{dropoff} = \left\{ \begin{array}{l} : ride_1 : dropoffTime "2013-01-01 15:18:10". \\ : ride_1 : tripttime 382. \end{array} \right\}. S_{fare} = \left\{ \begin{array}{l} : trans_1 : fare \ 7. \\ : trans_1 : pickupTime "2013-01-01 15:11:48". \end{array} \right\}.
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The query for continuous computation:
"hourly riding rate of active taxies of last 1000
payment transactions"

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"hourly riding rate of active taxies of last 1000 payment transactions"

Continuous Query in SPARQL-like language

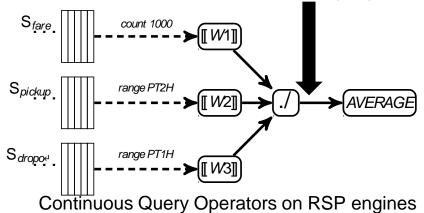
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On boosting the processing throughput of RSP engines



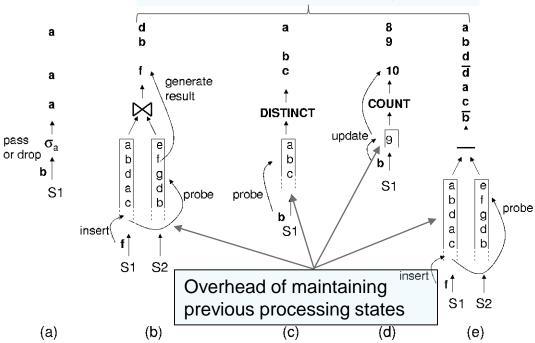
- Experience from Implementation CQELS Execution Framework: Using off-the-shelf data structures and algorithms are not enough!!??
 - ➤ Hardly can reach 10000 operator executions/second on large windows(100k-1M entries)
 - ➤ Big overhead of using row-based data structures
- Bottom-up perspective: investigating closely to data structures and algorithms
 - > Highly efficient data structures for maintaining processing states
 - > Sophisticate *incremental evaluation algorithms* of query operators



Incremental Evaluation for continuous operators in a Nutshell



stateful sliding window operators: reuse previous computing effort



L. Golab, M. T. Özsu, Data stream management, Synthesis Lectures on Data Management (2010) 1–73.





Incremental Evaluation of Continuous Queries

over RDF Stream: Issues and Challenges

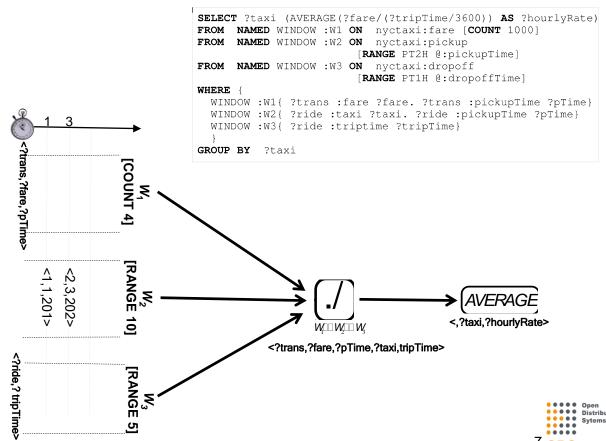


- Row-based data structure is not suitable for :
 - very small RDF data elements (encoded as fixed-size integers)
 - unusually large number individual data points (millions of mappings/RDF nodes are generated/evicted per second)
- Timestamping or negative-tuple solutions for incremental computation of RDF data elements and mappings have technical issues:
 - Auxiliary data (extra timestamps or negative tuples) might be bigger than original data
 - Other limitations of state-of-art techniques (double computation in evicting expired computing state)



Query: sliding data flow



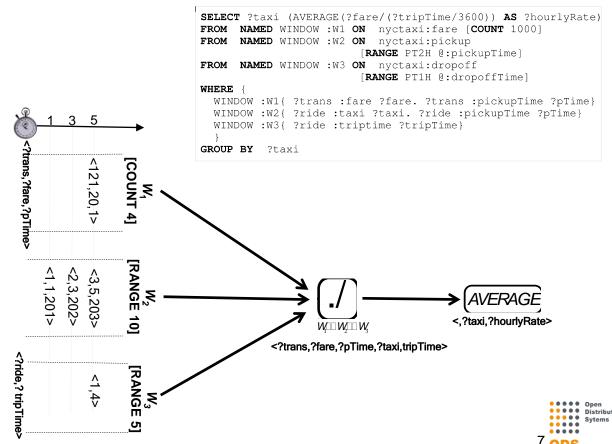


<?ride,?pTime,?taxi>



Query: sliding data flow



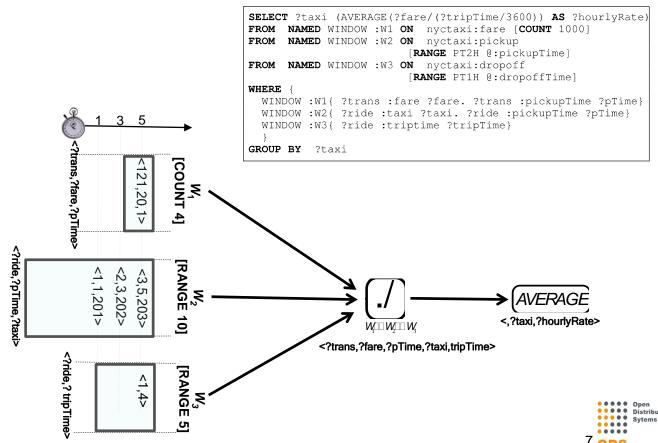


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Query: sliding data flow

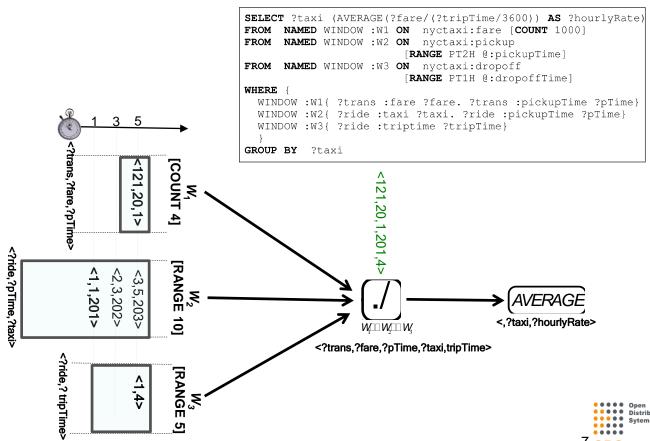






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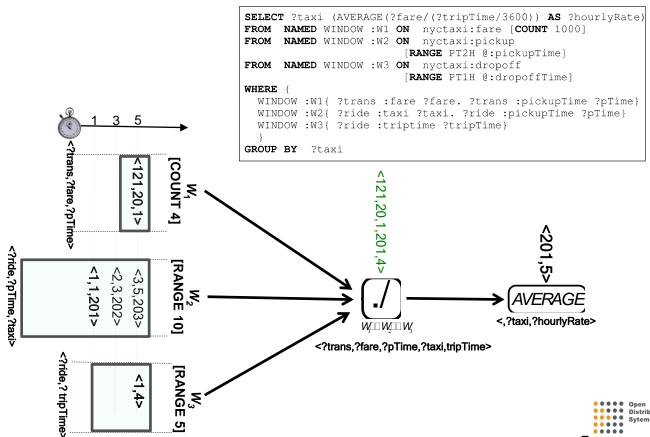






Query: sliding data flow

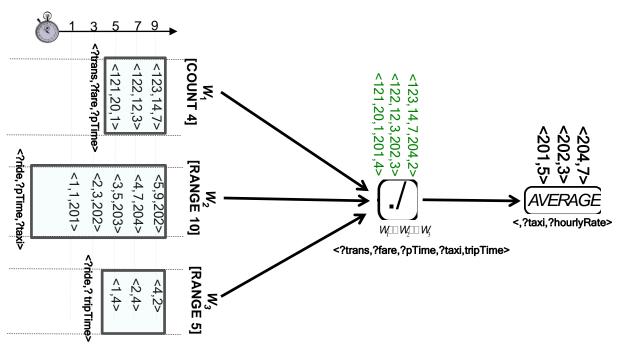






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Query: generating processing state

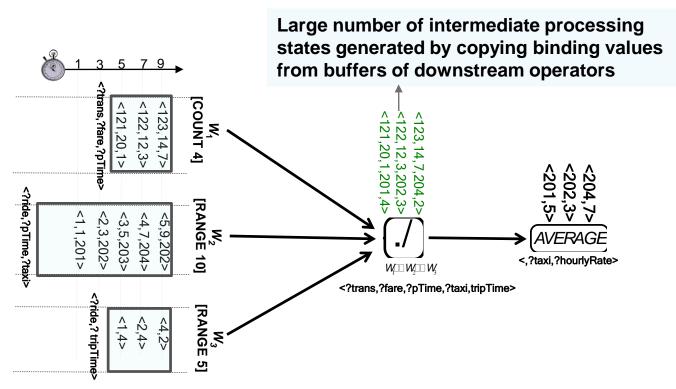








Query: generating processing state

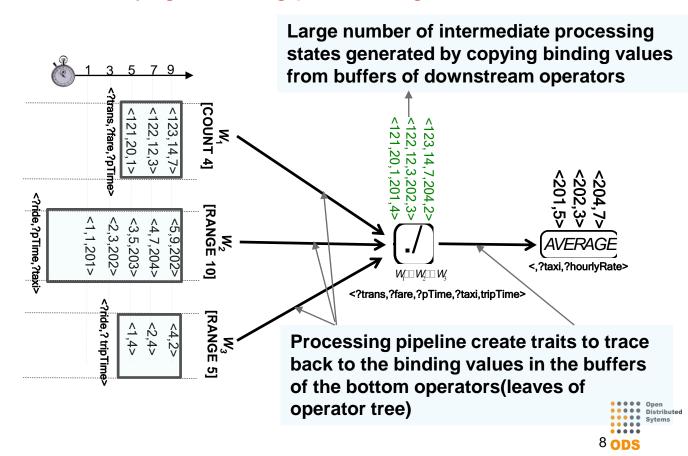






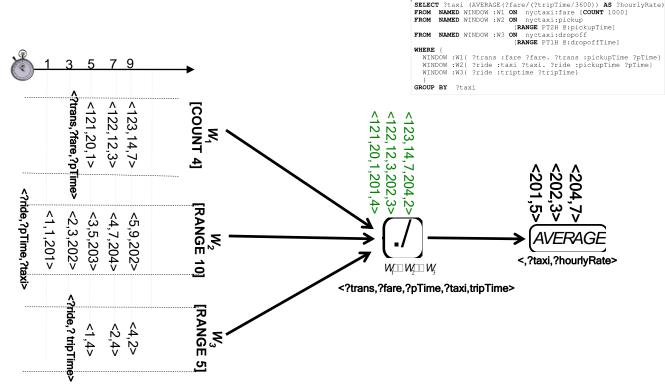


Query: generating processing state





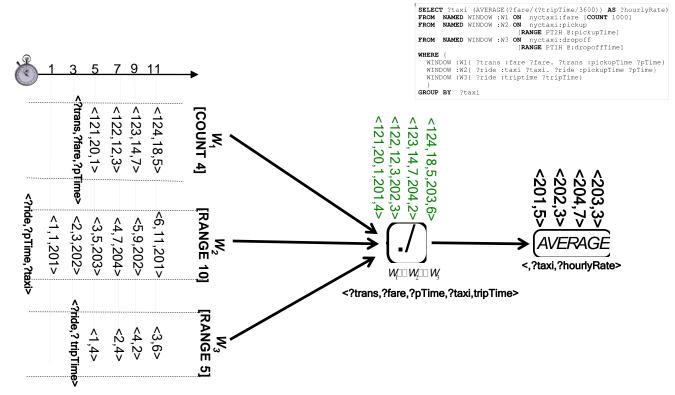








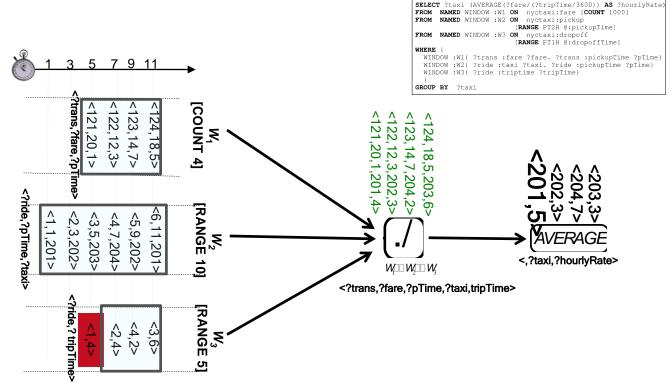








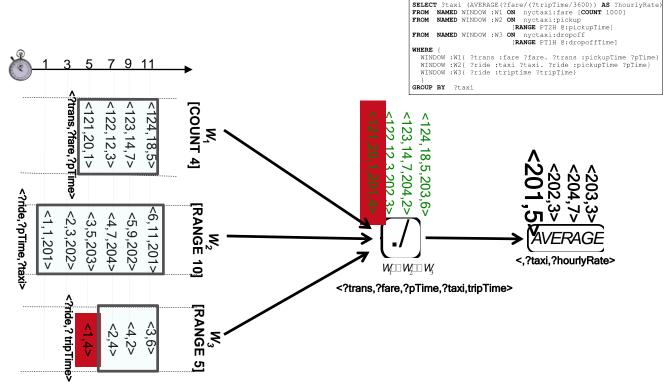








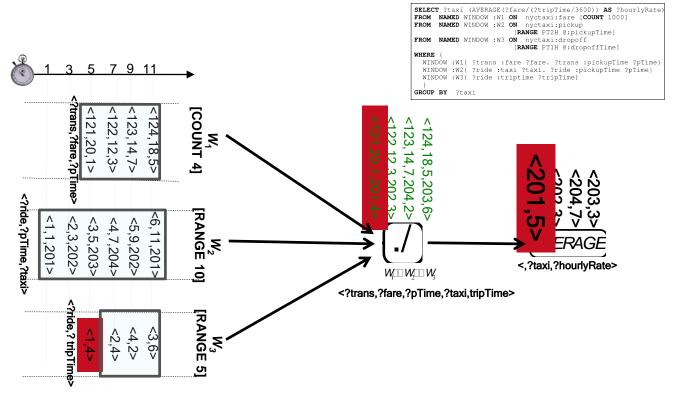








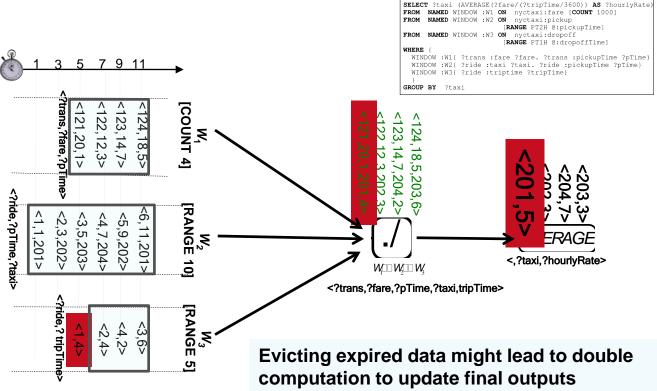
















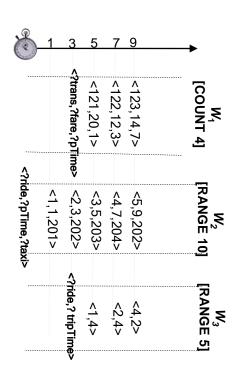




- Operator-aware data structures designed for:
 - Bookkeeping how the processing states were generated by the the query operators
 - Indexing windowing buffers tailored for query operators' behaviors
- Algorithms for incremental evaluations driven by operator-aware data structures



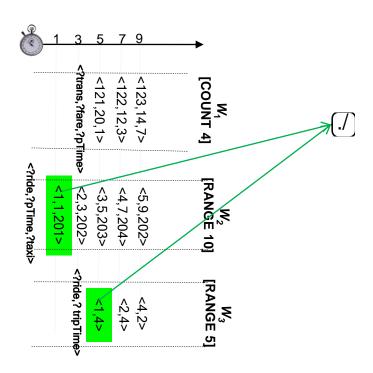








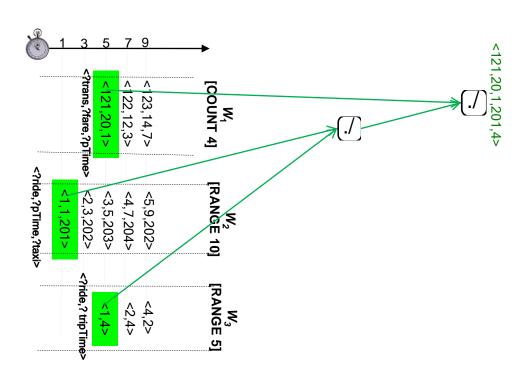








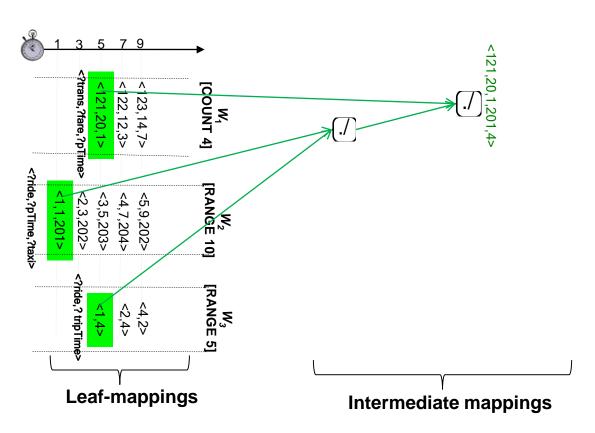








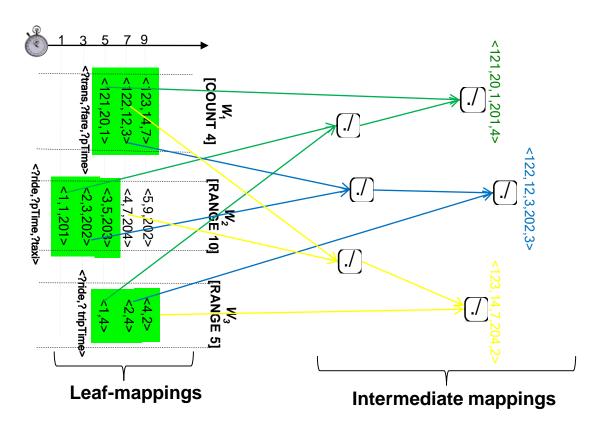








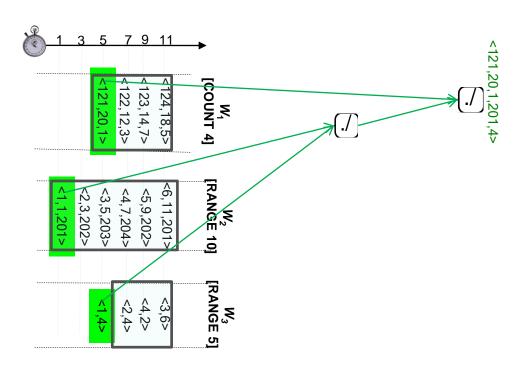








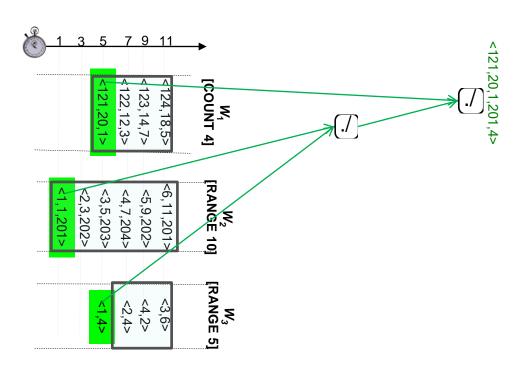








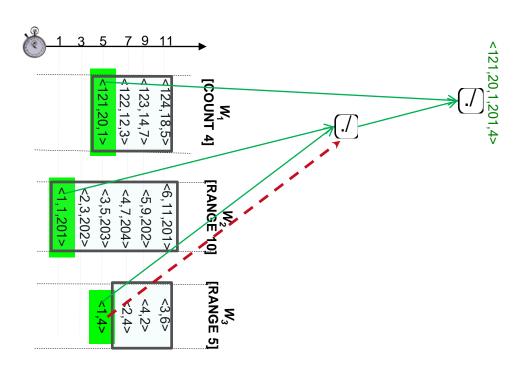








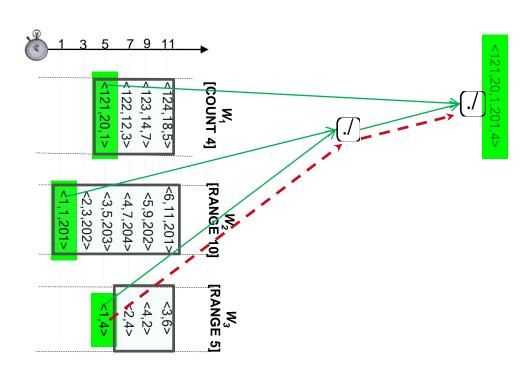












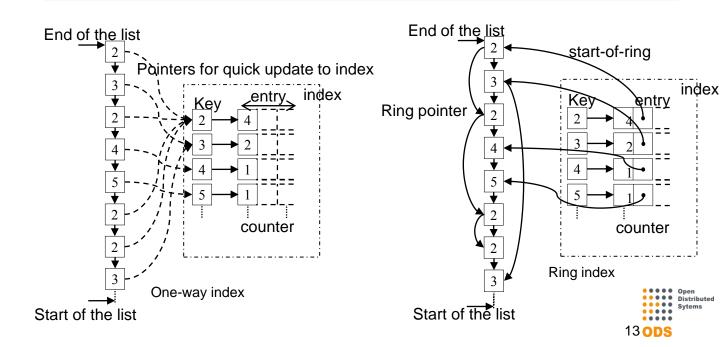




Ring Indexes on bags of mappings



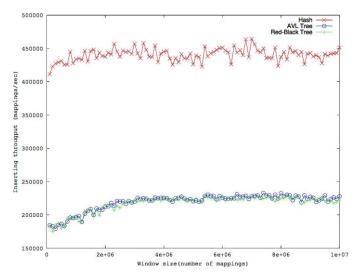
- Operator-aware indexes for quick lookup operations
- Low maintenance cost for fast insert/delete operations

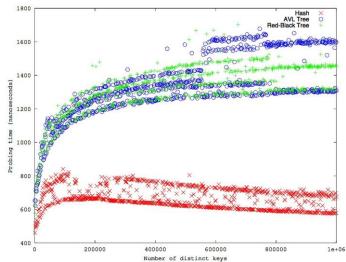




Throughputs of Ring indexes







Insert throughput for 1M keys

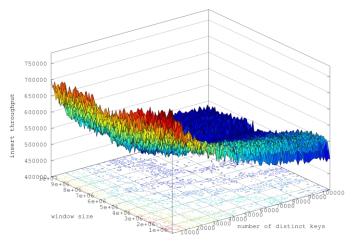
Probing time for 1M-mapping windows

Hash outperforms over AVL Tree and Red-Black Tree

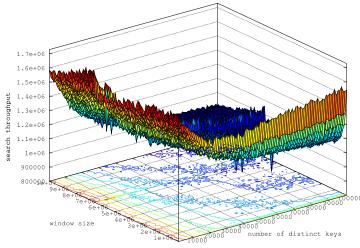


Throughputs of Ring indexes (Cont.)





Inserting throughput: 500-900k



Probing/searching throughput: 1M-1.6M

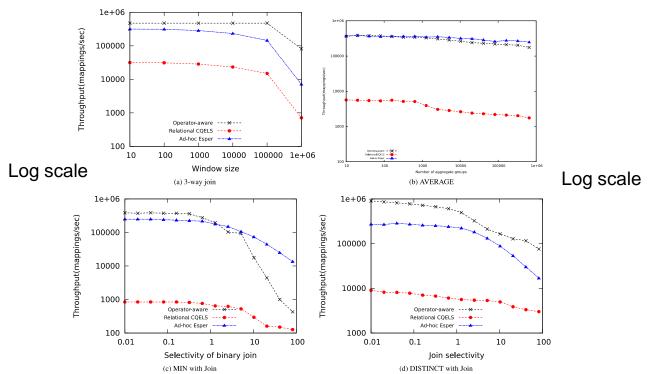




Throughputs of Query Operators



Sytems



- Operator-aware implementations outperform to relational implementations
- ... are marginally faster than ad-hoc implementations of ESPER in most cases



Throughputs & Memory footprint



| | | SRBer | nch (tripl | es/sec) | | LSBench (triples/sec) | | | | | | | | | |
|---------|-------|-------|------------|---------|----------|-----------------------|-------|-------|-------|-------|-------|----------|--|--|--|
| | Q_1 | Q_4 | Q_5 | Q_8 | Q_{10} | Q_1 | Q_2 | Q_3 | Q_4 | Q_5 | Q_6 | Q_{10} | | | |
| R-CQELS | 1214 | 820 | 47 | 1774 | 3343 | 24122 | 8462 | 9828 | 1304 | 7459 | 3491 | 2326 | | | |
| CQELS | 25147 | 20161 | 13966 | 22278 | 29463 | 118924 | 96789 | 88647 | 60467 | 52890 | 44391 | 103698 | | | |

Processing Throughputs: 5-12 times more than relational

| | Multiway Join (MB) | | | | | SRBench (MB) | | | | | | LSBench (MB) | | | | | | |
|---------|--------------------|-------|-------|-------|-------|--------------|-------|-------|-------|-----------------|-------------|--------------|-------|-------|-------|----------|--|--|
| | 2 | 3 | 4 | 6 | 8 | Q_1 | Q_4 | Q_5 | Q_8 | $Q_{10} \mid Q$ | Q_1 Q_2 | Q_3 | Q_4 | Q_5 | Q_6 | Q_{10} | | |
| R-CQELS | 25.61 | 28.67 | 38.95 | 49.24 | 54.61 | 457 | 745 | 834 | 620 | 488 38 | 35 404 | 420 | 490 | 502 | 560 | 420 | | |
| CQELS | | | | | | | | | | 218 3 | 4 370 | 380 | 398 | 389 | 402 | 370 | | |
| ESPER | 8.93 | 12.04 | 15.13 | 21.26 | 27.44 | | | | | · | | | | | | | | |

Memory Footprint: twice less memory than relational and 20-50% less than ESPER





Summary



- Incremental evaluation algorithms based on operator-aware data structures:
 - Overcome technical issues on traditional incremental evaluation techniques/algorithms
 - Perform several orders of magnitude faster than relation-based implementations
- Throughputs on operator-aware operations on processing state:
 - ❖Up to 1 million of updates/sec vs. 10k of relation-based one
 - ❖Up to 1.6 million lookup operations/second
 - Outperform over relational operations by order of magnitutes
 - ❖Consume twice less memory than relation-based implementations
- The implementation will be open sourced in the next release of CQELS(cqels.org)

