

Semantic Web Technology for Agent Communication Protocols

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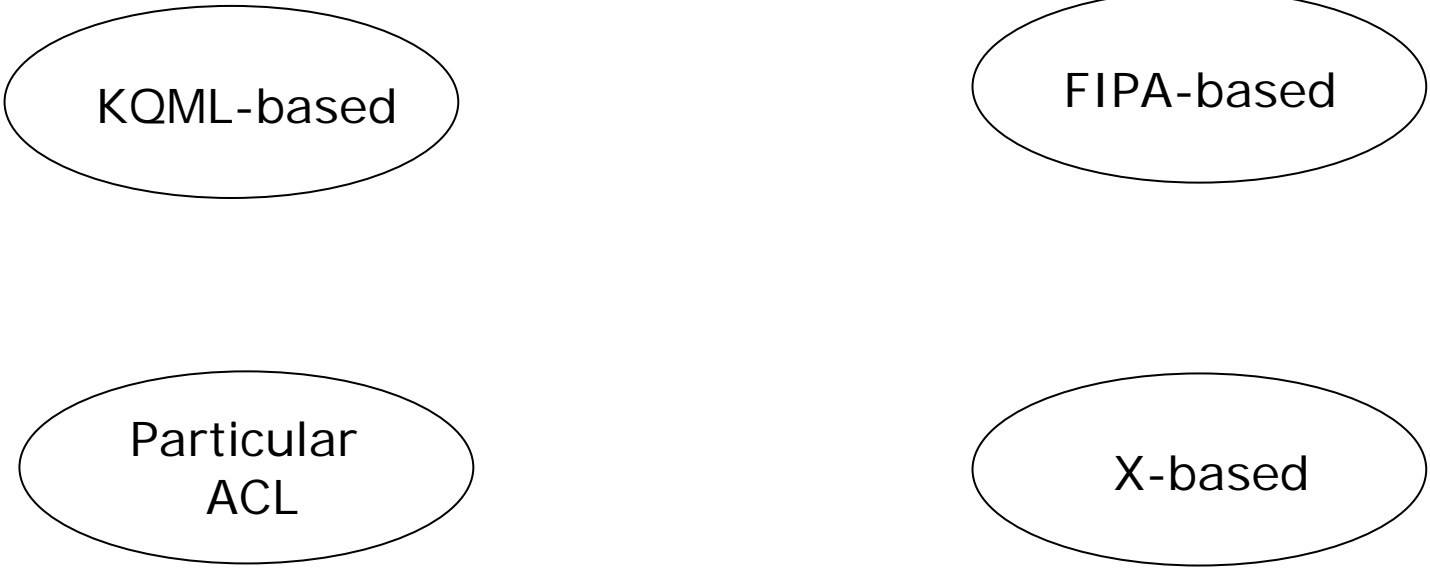


Outline

- Motivation and goal
- Ontology for communication acts
- Protocol descriptions
- Relationships between protocols
- Conclusions

Motivation

Right Now!



KQML-based

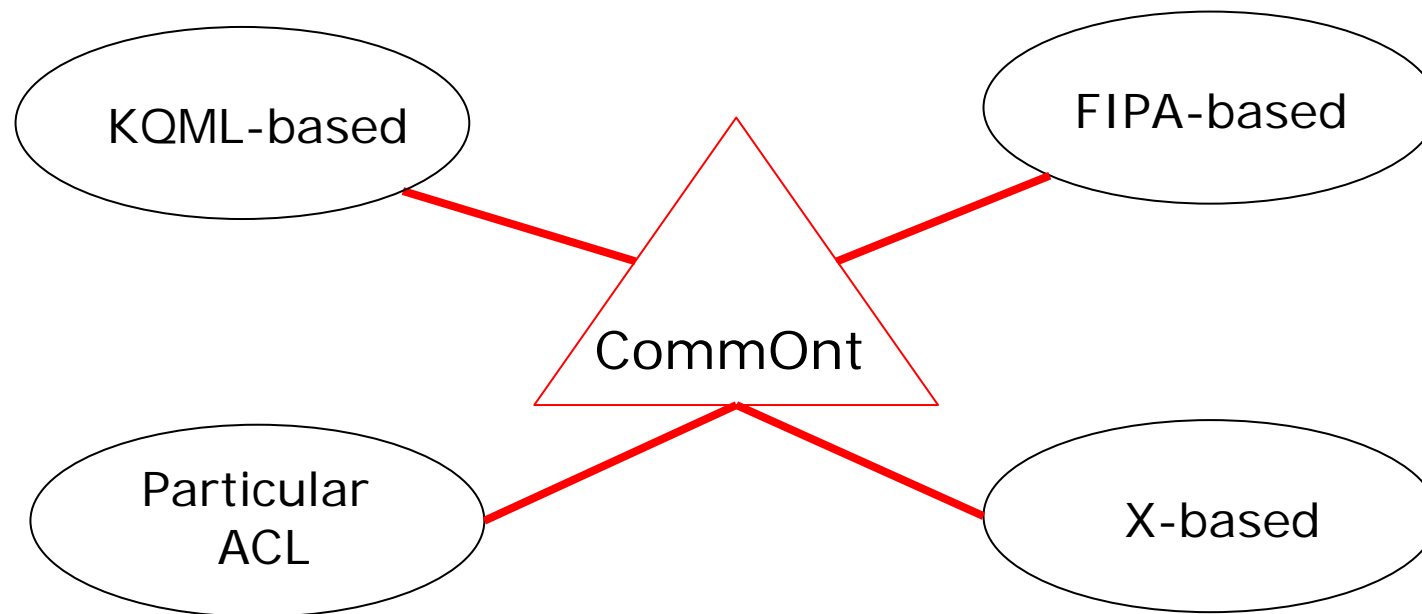
FIPA-based

Particular
ACL

X-based

Motivation

Where we want to arrive at!



Exchangeability of communication acts

Motivation

- Exchangeability of communication acts is not enough
- Sharing of communication protocols is needed

Proposal: Goal

- *Referential* representations for
 - communication acts
 - communication protocolsusing Semantic Web technology.

Proposal: Contributions

- To favour a flexible agent interoperation.
- To facilitate customization of communication protocols.
- A basis for reasoning about protocol relationships.
- Take account of semantics in protocol representations.

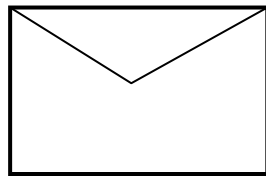
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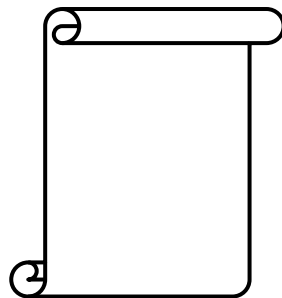
CommOnt design criteria(1)

■ Speech Acts Theory

- A communication act is basically composed of an envelope and a content



Intention + communication information



Object of the intention

CommOnt design criteria(2)

- Social commitments approach
 - Objective and verifiable semantics
 - Formalization:
 - Commitment
 - $C(x, y, p)$
 - Conditional commitment
 - $CC(x, y, c, p)$
 - $CC(x, y, c, p) \wedge c \rightarrow C(x, y, p)$

CommOnt design criteria(3)

- Axiomatization of communication acts with Event Calculus:
 - First-order theory for reasoning about actions
 - **Events** (actions) initiate and terminate fluents
 - **Fluents** are propositions whose value is subject to change over time

Event Calculus predicates

1. $Initiates(a, f, t)$ means that f holds after event a at time t .
2. $Terminates(a, f, t)$ means that f does not hold after event a at time t .
3. $Initially_P(f)$ means that f holds from time 0.
4. $Initially_N(f)$ means that f does not hold from time 0.
5. $Happens(a, t_1, t_2)$ means that event a starts at time t_1 and ends at t_2 .
6. $HoldsAt(f, t)$ means that f holds at time t .
7. $Clipped(t_1, f, t_2)$ means that f is terminated between t_1 and t_2 .
8. $Declipped(t_1, f, t_2)$ means that f is initiated between t_1 and t_2 .

From Pinar Yolum and Munindar P. Singh

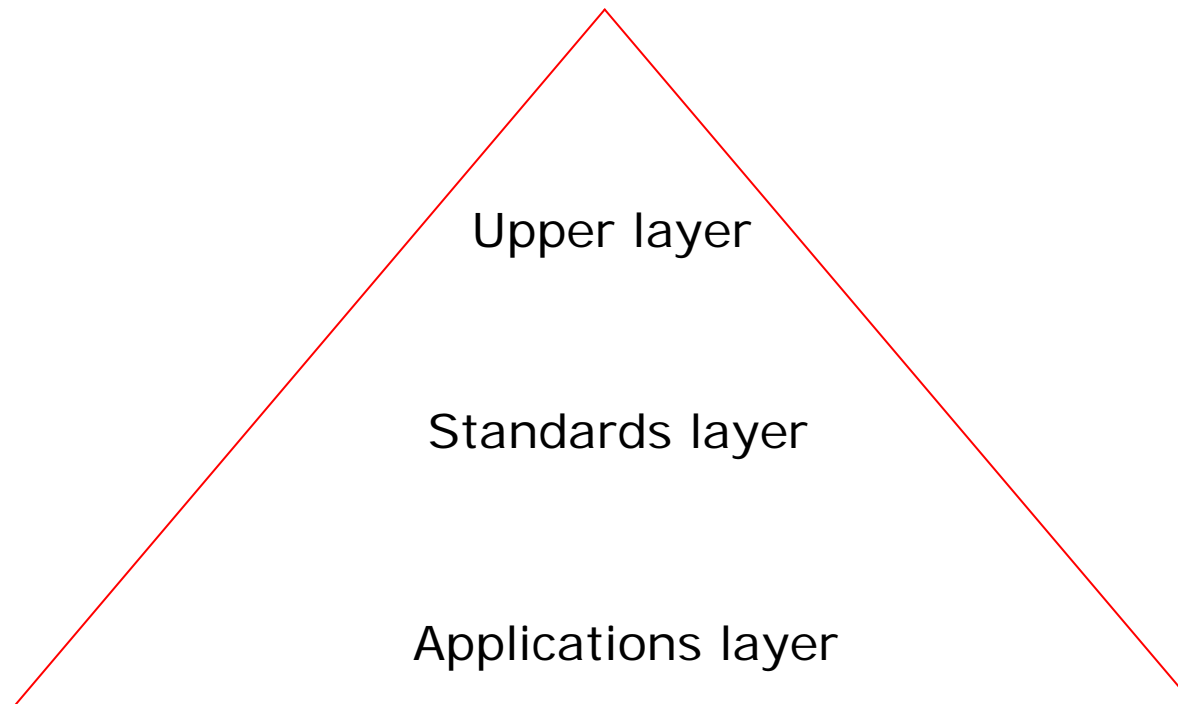
Rules

- To capture the dynamics of commitments

$$\textit{HoldsAt}(\textit{C}(x, y, p), t) \wedge \textit{Happens}(e(x), t) \wedge \textit{Initiates}(e(x), p, t) \rightarrow \\ \textit{Terminates}(e(x), \textit{C}(x, y, p), t).$$

CommOnt design criteria(4)

- Materialization of CommOnt: OWL ontology



CommOnt upper layer

`CommunicationAct` \sqsubseteq $\forall \text{hasSender.Actor} \sqcap =1.\text{hasSender} \sqcap$
 $\forall \text{hasReceiver.Actor} \sqcap$
 $\forall \text{hasContent.Content}$

Main subclasses:

Assertive, Directive, Commissive, Expressive and Declarative

Other subclasses:

`Inquiry` \sqsubseteq `Directive`
`Request` \sqsubseteq `Directive`
`Responsive` \sqsubseteq `Assertive`

CommOnt upper layer

`Assertive` \equiv `CommunicationAct` \sqcap
 \exists `hasContent.Proposition` \sqcap
 \exists `hasCommit.Commitment`

Initiates(*Assertive*(*s*, *r*, *P*), *C*(*s*, *r*, *P*), *t*)

`Directive` \equiv `CommunicationAct` \sqcap
 \exists `hasContent.Action` \sqcap
 \exists `hasCommit.ConditionalCommitment`

Initiates(*Directive*(*s*, *r*, *P*), *CC*(*r*, *s*, *accept*(*r*, *s*, *P*), *P*), *t*)

CommOnt standards layer

FIPA-Inform	\sqsubseteq	Assertive
FIPA-Confirm	\sqsubseteq	Assertive
FIPA-Disconfirm	\sqsubseteq	Assertive
FIPA-Request	\sqsubseteq	Directive

KQML-Tell	\sqsubseteq	Assertive
KQML-Ask-If	\sqsubseteq	Inquiry \sqcap $\forall \text{hasContent.}(\forall \text{hasQuery. Proposition})$

- Very important for the interoperability goal:

FIPA-Inform	\equiv	KQML-Tell
KQML-Achieve	\sqsubseteq	FIPA-Request
KQML-Ask-If	\equiv	FIPA-Query-If
KQML-Achieve	\equiv	FIPA-Request $\sqcap \exists \text{hasContent.Achieve}$

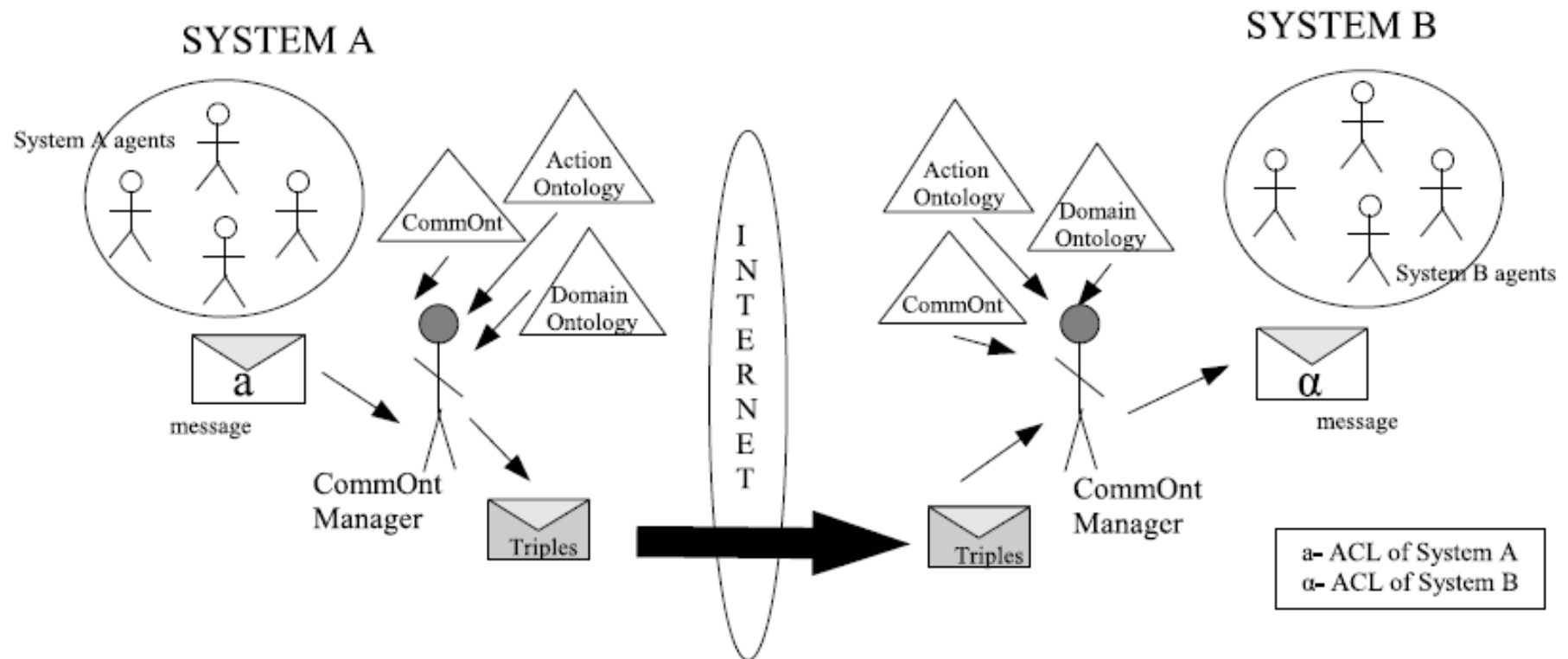
CommOnt applications layer

MedicineModify \equiv Request \sqcap =1.hasContent \sqcap
 \forall hasContent.(Overwrite \sqcap \exists hasSubject.Medicine)

LocationQuery \equiv Inquiry \sqcap =1.hasContent \sqcap
 \forall hasContent.(
 \forall hasQuery.(RefExpression \sqcap \exists hasSubject.Location))

VitalSignInform \equiv Responsive \sqcap =1.hasContent \sqcap
 \forall hasContent.(Proposition \sqcap \exists hasSubject.VitalSignData)

Communication process



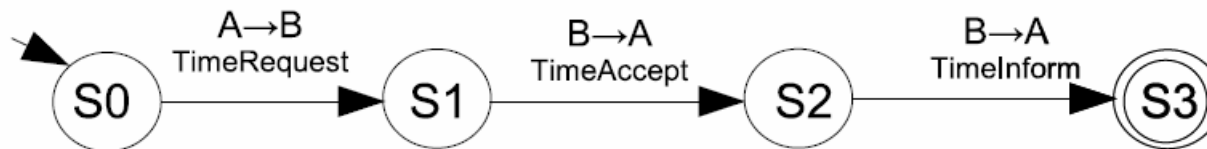
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Protocol descriptions

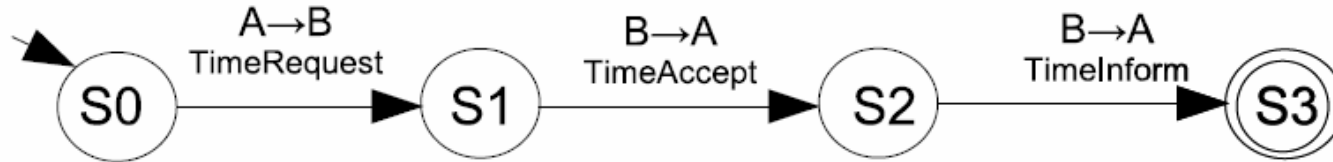
- Model: State Transition System
 - Transitions labeled with classes of Communication Acts
 - States associated to sets of fluents

Protocol AskTime



Protocol description

Protocol AskTime



$\text{Asktime} \equiv \text{Protocol} \sqcap \exists \text{hasInitialState.S0}$

$\text{S0} \equiv \text{State} \sqcap \exists \text{hasTransition.T01} \sqcap \exists \text{hasFluent.F0}$

$\text{S1} \equiv \text{State} \sqcap \exists \text{hasTransition.T12} \sqcap \exists \text{hasFluent.F1}$

$\text{S2} \equiv \text{State} \sqcap \exists \text{hasTransition.T23} \sqcap \exists \text{hasFluent.F2}$

$\text{S3} \equiv \text{FinalState} \sqcap \exists \text{hasFluent.F3}$

$\text{T01} \equiv \text{Transition} \sqcap \exists \text{hasCommAct.TimeRequest} \sqcap \exists \text{hasNextState.S1}$

$\text{T12} \equiv \text{Transition} \sqcap \exists \text{hasCommAct.TimeAccept} \sqcap \exists \text{hasNextState.S2}$

$\text{T23} \equiv \text{Transition} \sqcap \exists \text{hasCommAct.TimeInform} \sqcap \exists \text{hasNextState.S3}$

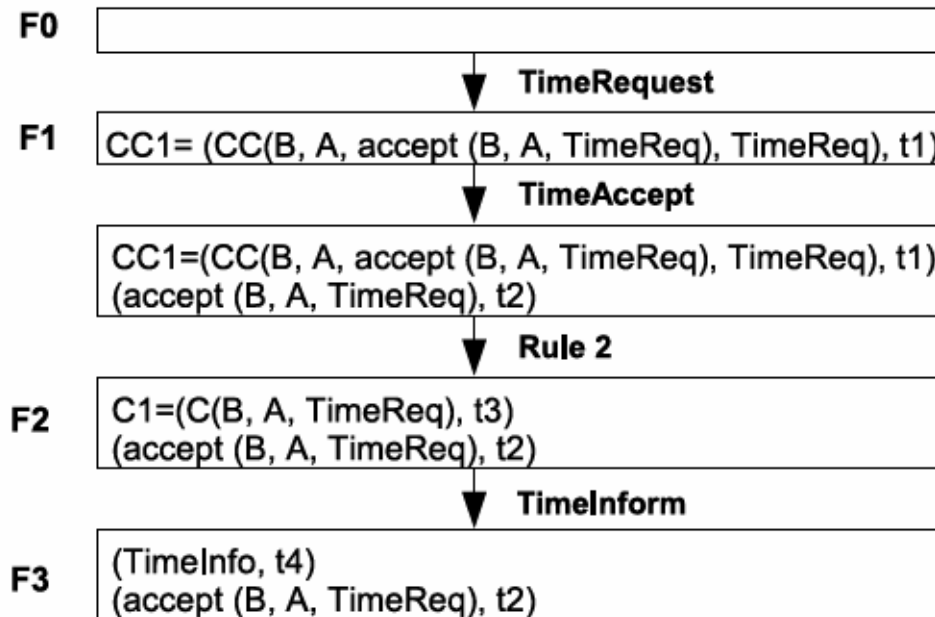
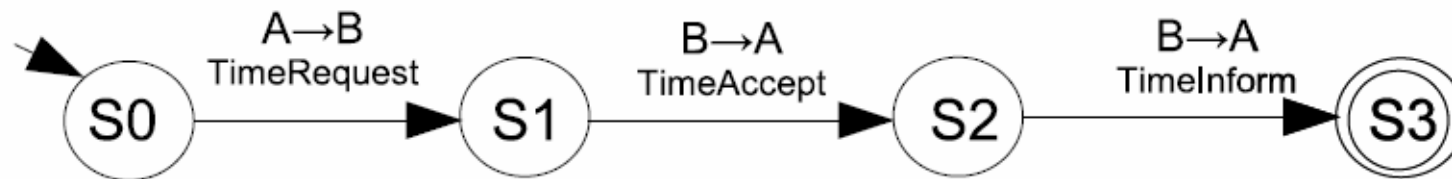
$\text{TimeRequest} \equiv \text{Request} \sqcap =1 \text{ hasContent.TimeReq}$

$\text{TimeAccept} \equiv \text{Accept} \sqcap =1 \text{ hasContent.TimeReq}$

$\text{TimeInform} \equiv \text{Responsive} \sqcap =1 \text{ hasContent.TimeInfo} \sqcap =1 \text{ inReplyTo.TimeRequest}$

Simulation of a protocol run

Protocol AskTime

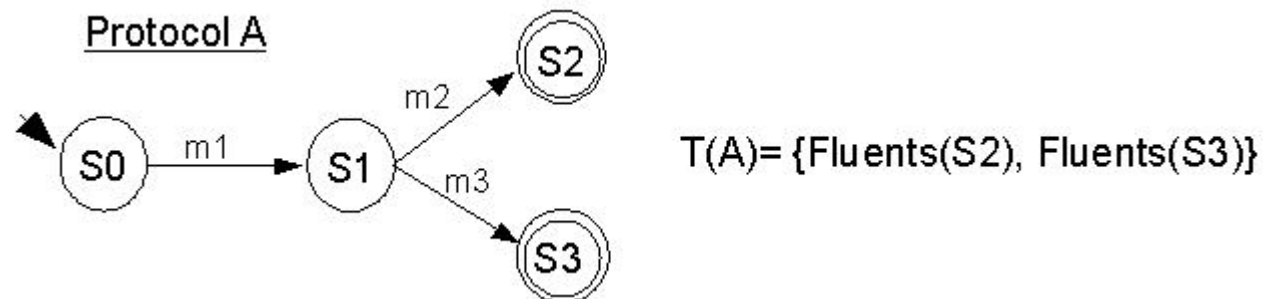


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Protocol relationships (1)

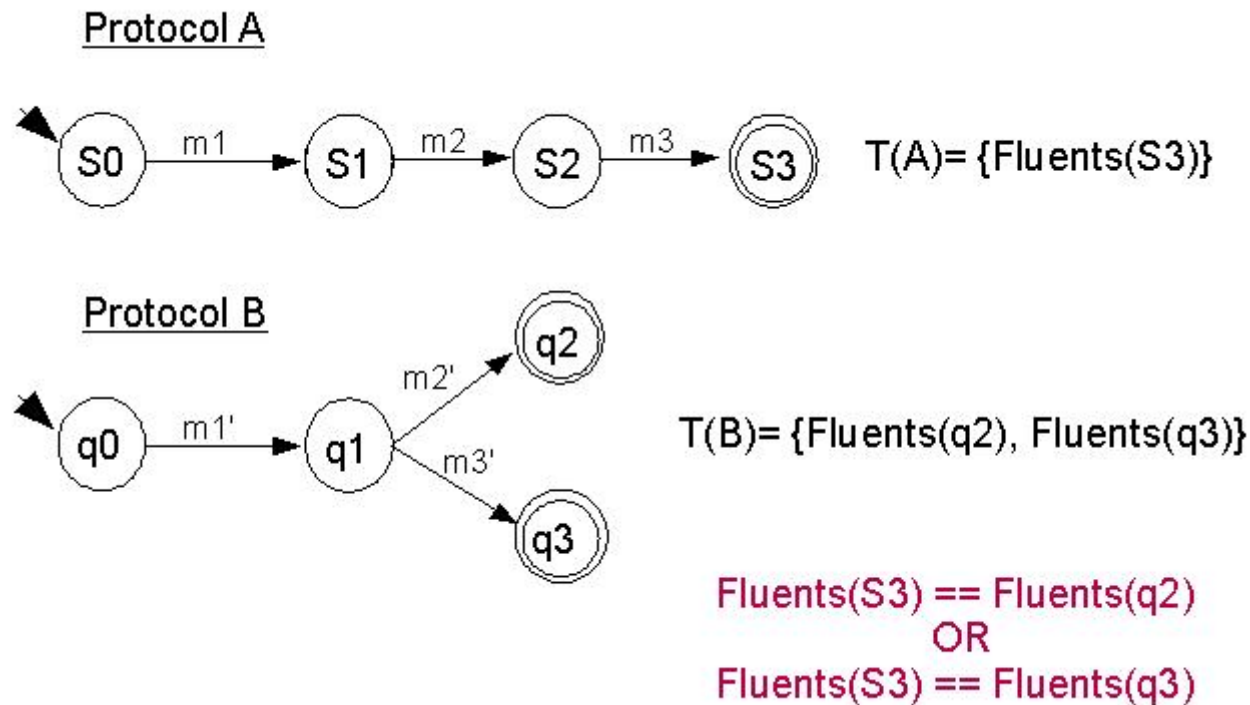
- $T(A)$ is the set of traces of A



- Protocol A is ***equivalent*** to protocol B if
“ $T(A) = T(B)$ ”

Protocol relationships (2)

- Protocol A is a **restriction** of protocol B if
“ $T(A) \subset T(B)$ ”



Protocol relationships (3)

- Protocol A is **specialized-equivalent** to protocol B if “ $T(A) \ll T(B)$ ”

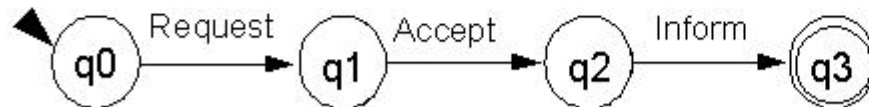
Protocol A



$T(A) = \{\text{Fluents}(S3)\}$

accept(r,s,TimeReq)
TimeInfo

Protocol B

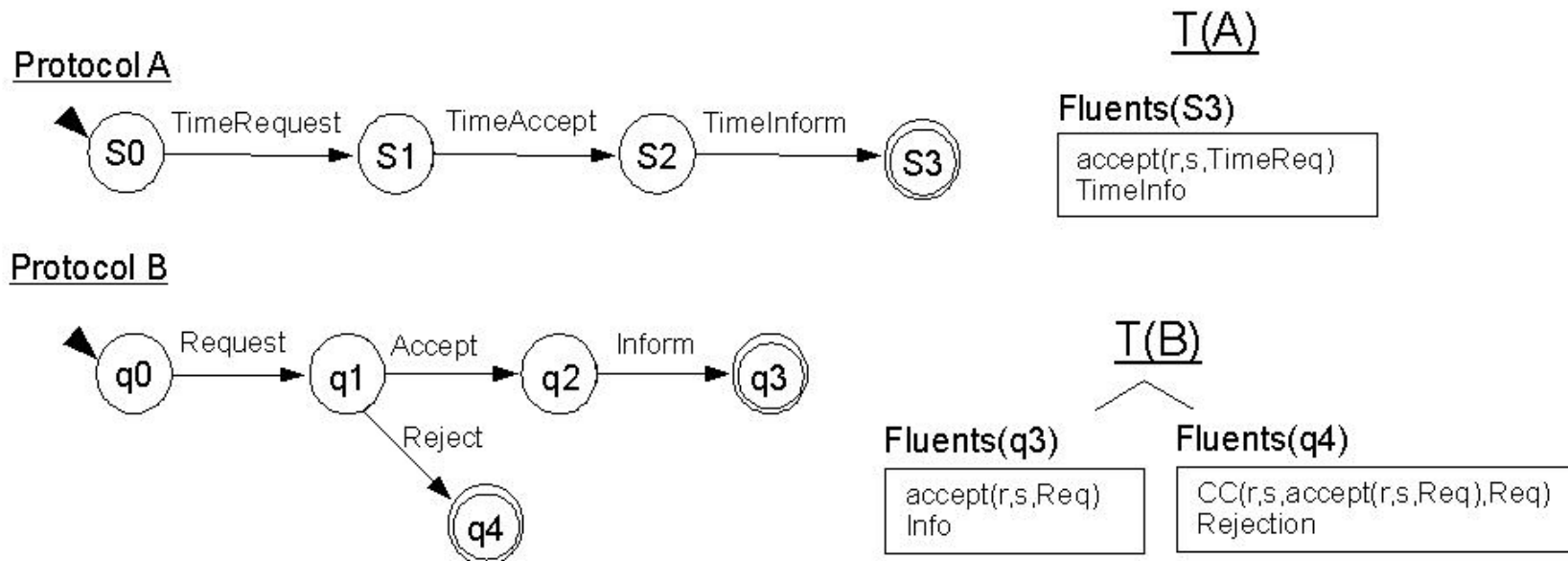


$T(B) = \{\text{Fluents}(q3)\}$

accept(r,s,Req)
Info

Protocol relationships (4)

- Protocol A is **specialized-restriction** to protocol B if “ $T(A) \subset \ll T(B)$ ”



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Conclusions

- Our proposal facilitates:
 - Using CommOnt ontology:
 - Management of semantic aspects when dealing with agent communication protocols
 - Customization of standard communication protocols
 - Support for discovering different kinds of relationships between protocols
 - Use of standard Semantic Web tools

Thanks for your attention!

