Algebras of ontology alignment relations

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Algebra of alignment relations

Benefits

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Ontology alignments



Ontology alignments



Ontology alignments



An alignment is a set of correspondences:

Definition (Correspondence)

Given two ontologies o and o' with associated entity languages Q_L and $Q_{L'}$, a set of alignment relations Θ , and a confidence structure over Ξ , a correspondence is a quadruple: $\langle e, e', r, n \rangle$, such that

•
$$e \in Q_L(o)$$
 and $e' \in Q_{L'}(o')$;

•
$$r \in \Theta$$
;

We will first restrict ourselves to correspondences as $\langle e, e', r \rangle$.

Stadt ? Town

Stadt ? Town

merging independent positive alignments:

Stadtgebiet \emptyset Municipality Stadtgebiet \leq Municipality

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What is Θ?

The matching process can be achieved by aggregating:

- independent sources of matching along different dimensions (based on instances rather than definitions);
- competing sources of matching (providing different opinions on the same correspondences);
- indirect sources of matching (using an intermediate ontology and composing the alignments).

This is becoming more common with the availability of alignments on the web.

How to combine these alignments?

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Definition (Relation algrebra (Tarski 1941))

An algebra of binary relations is a structure $\langle \Theta, \wedge, \vee, *, 1, 0, 1', \neg \rangle$ such that

- $\langle \Theta, \wedge, \vee, 1, 0 \rangle$ is a Boolean algebra;
- ▶ * is an associative internal composition law with (left and right) unity element 1', that distributes over ∨;
- $\blacktriangleright \neg$ is an internal involutive unary operator, that distributes over \lor, \land and $\ast.$

$$\begin{array}{l} ((\leq * \leq) \land \geq) \\ = (\leq \land \geq) \\ == \\ \leq \lor \geq ? \end{array}$$

 Γ is a set of jointly exhaustive and pairwise disjoint (JEPD) relations between two entities.

Definition (Powerset relation algrebra)

A powerset algebra of binary relations is a structure $\langle 2^{\Gamma},\cap,\cup,\cdot,\Gamma,\varnothing,\{=\},\cdot^{-1}\rangle$ such that

- $\langle 2^{\Gamma}, \cap, \cup, \Gamma, \varnothing \rangle$ is a Boolean algebra;
- ► · is an associative internal composition law with (left and right) unity element {=}, that distributes over ∪;
- \blacktriangleright \cdot^{-1} is an internal involutive unary operator, that distributes over $\cup,\,\cap\,$ and $\cdot.$

$$((\{<,=\} \cdot (\{<,=\}) \cap (\{>,=\})) = (\{<,=\}) \cap (\{>,=\})) = \{=\}$$
$$(\{<,=\} \cup \{>,=\}) = \{<,>,=\}$$

Because of the usual semantics of classes in ontology languages, it is possible to create a simple algebra corresponding to relations between sets:

- = classes are equivalent;
- $<\,$ the first class is subsumed by the second one;
- > the first class subsumes the second one;
- $\ensuremath{\emptyset}$ classes overlap but none subsumes the other;
- \perp classes are disjoint.

Lattice of relations among classes (or properties)



Lattice of relations among classes (or properties)



Using alignment algebra is not restricted to these very classical types of relations:

Adding instances:

- relation between instances: =, \perp ;
- ▶ relations between instances and classes: \in , \notin ;
- ▶ relations between classes and instances: \ni , $\not\ni$.

Adding other relations:

- part-of relations
- application relations

But beware of the JEPD constraint: these are obtained by product, not addition.

Motivation: disjunctive alignments

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Stadt ? Town

Stadt $\{=,<,>,\emptyset\}$ Town

Stadt
$$\{=,<,>,\emptyset\}$$
 Town

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Stadtgebiet () Municipality Stadtgebiet \leq Municipality

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Stadtgebiet
$$\{=, <, \emptyset\}$$
 Municipality

Now we have several choices for aggregating alignments:

- disjunction is natural for combining independent alignments;
- conjunction is useful for aggregating competing alignments.

Alignment disjunction example

$$\mathbf{A_8} = A_6 \lor A_4 : o \to o'$$

 ${\sf Stadt}\{=,<,>,\check{\mathbb{Q}}\}{\sf Town}$ ${\sf Stadtgebiet}\{\check{\mathbb{Q}},\bot\}{\sf Municipality}$

Alignment disjunction example

 $\begin{array}{ll} \textbf{A_6}: & o \rightarrow o'\\ \text{Konstruktion}\{\bot\}\text{Municipality}\\ & \text{Stadt}\{<\}\text{Town}\\ \text{Stadtgebiet}\{\emptyset\}\text{Municipality} \end{array}$

 $A_4: o \rightarrow o'$

$$\label{eq:Konstruktion} \begin{split} & \mathsf{Konstruktion}\{=,<,>,\emptyset,\bot\} \\ & \mathsf{Municipality} \\ & \mathsf{Stadt}\{=,<,>,\emptyset\} \\ & \mathsf{Town} \\ & \mathsf{Stadtgebiet}\{\bot\} \\ & \mathsf{Municipality} \end{split}$$

$$\begin{split} \mathbf{A_8} &= A_6 \lor A_4: \ o \to o' \\ & \mathsf{Konstruktion}\{=,<,>, \Diamond, \bot\} \mathsf{Municipality} \\ & \mathsf{Stadt}\{=,<,>, \check{\Diamond}\} \mathsf{Town} \\ & \mathsf{Stadtgebiet}\{\check{\Diamond}, \bot\} \mathsf{Municipality} \end{split}$$

$$\textbf{A_3}: o \rightarrow o'$$

 $\mathsf{Konstruktion}\{\bot\}\mathsf{Municipality}$

$$A_5: o \rightarrow o'$$

Stadtgebiet $\{>, \emptyset\}$ Municipality

 ${\sf Stadt}\{<\}{\sf Town}$ ${\sf Stadtgebiet}\{\bot, \emptyset\}{\sf Municipality}$

$$\begin{split} \mathbf{A_6} &= A_3 \wedge A_5: \ o \to o' \\ & \text{Konstruktion}\{\bot\} \text{Municipality} \\ & \text{Stadt}\{<\} \text{Town} \\ & \text{Stadtgebiet}\{\emptyset\} \text{Municipality} \end{split}$$

Alignment conjunction example

$$A_3: o \rightarrow o'$$

$$\label{eq:Konstruktion} \begin{split} & \mathsf{Konstruktion}\{\bot\} \mathsf{Municipality} \\ & \mathsf{Stadt}\{=,<,>,\emptyset,\bot\} \mathsf{Town} \\ & \mathsf{Stadtgebiet}\{>,\emptyset\} \mathsf{Municipality} \end{split}$$

 $\begin{array}{lll} \textbf{A_5}: & o \to o' \\ & \text{Konstruktion}\{=,<,>, \Diamond, \bot\} \\ & \text{Municipality} \\ & \text{Stadt}\{<\} \\ & \text{Town} \end{array}$

$$\mathsf{Stadtgebiet}\{\bot, \emptyset\}\mathsf{Municipality}$$

$$\begin{split} \mathbf{A_6} &= A_3 \wedge A_5: \ o \to o' \\ & \text{Konstruktion}\{\bot\} \text{Municipality} \\ & \text{Stadt}\{<\} \text{Town} \\ & \text{Stadtgebiet}\{\emptyset\} \text{Municipality} \end{split}$$

Alignment composition is very useful (and algebra of relations are very natural):

 when one has no direct alignments between two ontologies but intermediate alignments;



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Alignment composition is very useful (and algebra of relations are very natural):

- when one has no direct alignments between two ontologies but intermediate alignments;
- ▶ when one wants to saturate a network of ontologies and alignments.



Composition table

$$\{<, \emptyset\} \cdot \{<, =\}$$

= $\{<\} \cdot \{<\} \cup \{<\} \cdot \{=\} \cup \{\emptyset\} \cdot \{<\} \cup \{\emptyset\} \cdot \{=\}$
= $\{<\} \cup \{<\} \cup \{<, \emptyset\} \cup \{\emptyset\} \cup \{\emptyset\}$
= $\{<, \emptyset\}$

 $\mathsf{A}\{<,\emptyset\}\mathsf{B}\{<,=\}\mathsf{C}\Rightarrow\mathsf{A}\{<,\emptyset\}\mathsf{C}$

$$A_1: o \rightarrow o''$$

Konstruktion $\{\bot\}$ Commune Stadtgebiet $\{>\}$ Ville

$$A_2: o'' \rightarrow o'$$

$$\label{eq:commune} \begin{split} \mathsf{Commune}\{>,=\} \mathsf{Municipality}\\ \mathsf{Ville}\{\emptyset\} \mathsf{Municipality} \end{split}$$

$$f A_3 = A_1 \cdot A_2 : o o o'$$

Konstruktion $\{\bot\} \cdot \{>, =\}$ Municipality
Stadtgebiet $\{>\} \cdot \{\emptyset\}$ Municipality

$$\label{eq:A1} \mathsf{A}_1: \ o \to o''$$
 Konstruktion $\{\bot\}$ Commune

 $Stadtgebiet {>} Ville$

$$\mathbf{A_2}: o'' \rightarrow o'$$

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$$\begin{split} \mathbf{A}_3 &= A_1 \cdot A_2 : \ o \to o' \\ & \text{Konstruktion} \{\bot\} \cdot \{>, =\} \cap \{\bot\} \cdot \{\emptyset\} \\ & \text{Municipality} \\ & \text{Stadtgebiet} \{>\} \cdot \{\emptyset\} \\ & \text{Municipality} \end{split}$$

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$$\begin{split} \mathbf{A}_3 &= A_1 \cdot A_2 : \ o \to o' \\ & \text{Konstruktion} \{\bot\} \cap \{<, \emptyset, \bot\} \\ & \text{Municipality} \\ & \text{Stadtgebiet} \{>, \emptyset\} \\ & \text{Municipality} \end{split}$$

$$\begin{array}{lll} \textbf{A_1}: & o \to o'' \\ \text{Konstruktion} \{\bot\} \text{Commune} \\ & \text{Stadtgebiet} \{>\} \text{Ville} \\ & \text{Konstruktion} \{\bot\} \text{Ville} \end{array}$$

$$A_2: o'' \rightarrow o'$$

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The algebra of relations between the same set of entities is not absolute. There can be many different set of base relations which depends on the granularity or the point of view of the representation.



Problem: can we deal with alignments at different granularities?

$$f A_7[\{ot,ot\}]: o o o'$$

Stadt $\{ot\}$ Town
Stadtgebiet $\{ot\}$ Municipality

$$\pmb{\mathsf{A_5}:} \ o \to o'$$

 ${\sf Stadt}\{<\}{\sf Town}$ ${\sf Stadtgebiet}\{\bot, \check{Q}\}{\sf Municipality}$

$$\begin{array}{ll} \textbf{A}_4 = \Uparrow A_7: \ o \to o' & \textbf{A}_9[\{\bot, \swarrow\}] = \Downarrow A_5 \\ \text{Stadt}\{=, <, >, \emptyset\} \text{Town} & \text{Stadt}\{\swarrow\} \text{Town} \\ \text{Stadtgebiet}\{\bot\} \text{Municipality} & \text{Stadtgebiet}\{\bot, \swarrow\} \text{Municipality} \end{array}$$

Thanks to the use of compatible algebras.



Definition (Confidence structure)

A confidence structure is an ordered set of degrees $\langle \Xi, \leq \rangle$ for which there exists a greatest element \top and a smallest element \perp .

A correspondence $\langle e, e', r, n \rangle$ entails another correspondence $\langle e, e', r', n' \rangle$ if and only if $r \subseteq r'$ and $n \ge n'$.

 $\mathsf{Stadt}\{<\}_{.9}\mathsf{Town}\models\mathsf{Stadt}\{<,=\}_{.8}\mathsf{Town}$

So the order between correspondances is the order induced by \subseteq, \geq .

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 \models Stadt Γ_{\top} Town

So the order between correspondances is the order induced by \subseteq, \geq .

Logical operator with confidence example

$$\label{eq:stadt} \begin{split} \mathsf{Stadt} \{=,<,\bot\}_{.2} \mathsf{Town} \\ \mathsf{Stadtgebiet} \{\bot, \emptyset\}_{.1} \mathsf{Municipality} \end{split}$$

 $\wedge = \cap^{\mathsf{max}}$ yields:

$$\label{eq:stadt} \begin{split} & \mathsf{Stadt}\{=\}_{1.}\mathsf{Town} \\ & \mathsf{Stadtgebiet}\{\bot\}_{.2}\mathsf{Municipality} \end{split}$$

Logical operator with confidence example (relaxed)

$$\begin{split} \mathsf{Stadt}\{=,<\}_{.5}\mathsf{Town}\\ \mathsf{Stadt}\{=,\bot\}_{.2}\mathsf{Town}\\ \mathsf{Stadtgebiet}\{\bot\}_{.1}\mathsf{Municipality}\\ \lor = \cap^{\mathsf{min}} \text{ yields:} \end{split}$$

 $\mathsf{Stadt}\{=\}_{1.}\mathsf{Town}$

 $\mathsf{Stadtgebiet}\{\bot, \emptyset\}_{.2}\mathsf{Municipality}$

$$\label{eq:stadt} \begin{split} & \mathsf{Stadt}\{=\}_{1.}\mathsf{Town}\\ & \mathsf{Stadtgebiet}\{\bot, \emptyset\}_{.2}\mathsf{Municipality}\\ & \mathsf{Stadtgebiet}\{\bot\}_{.1}\mathsf{Municipality} \end{split}$$

 $\wedge = \cup^{weightedsum}$ yields:

 $\label{eq:stadt} \begin{array}{l} \mathsf{Stadt}\{<,=\}_{.83}\mathsf{Town}\\ \mathsf{Stadt}\{\bot,=\}_{.4}\mathsf{Town}\\ \mathsf{Stadtgebiet}\{\bot,\check{\Diamond}\}_{.16}\mathsf{Municipality}\\ \end{array}$

- Alignment relation algebras are an adequate tool for representing disjunctions of relations;
- Moreover, they are naturally useful for many operation: aggregation, composition, granularity;
- They can be made compatible with confidence measures;
- They usually provide a better way to compare alignments;

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- Moreover, they are naturally useful for many operation: aggregation, composition, granularity;
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- They usually provide a better way to compare alignments;

- Made for using in alignment management;
- Not implemented yet feature (likely in Alignment API 4.0).

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