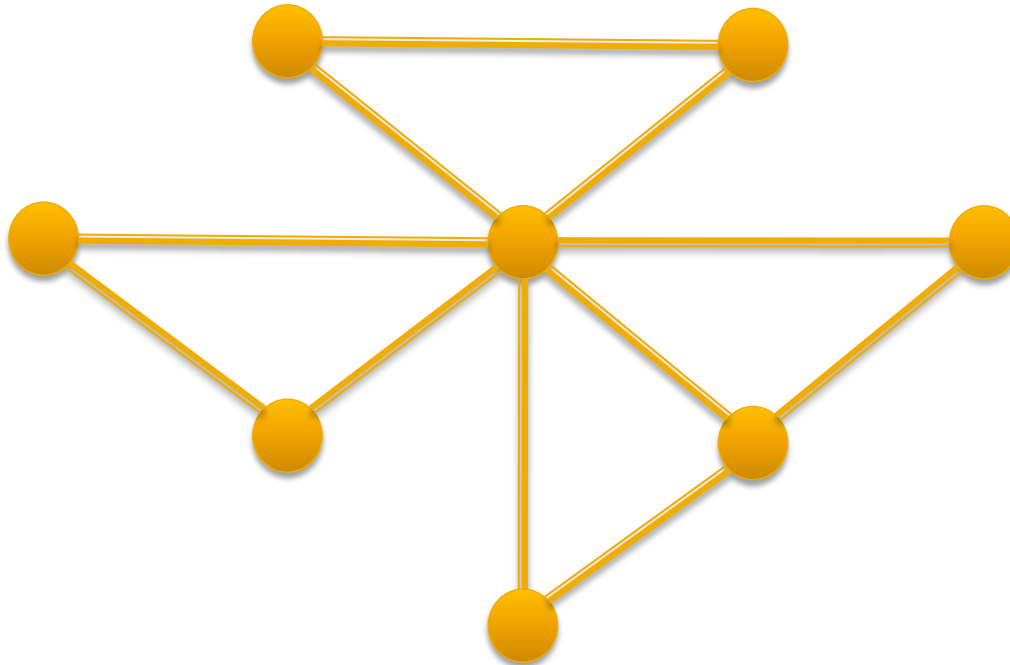


Andreas Björklund

The Path Taken for k-Path

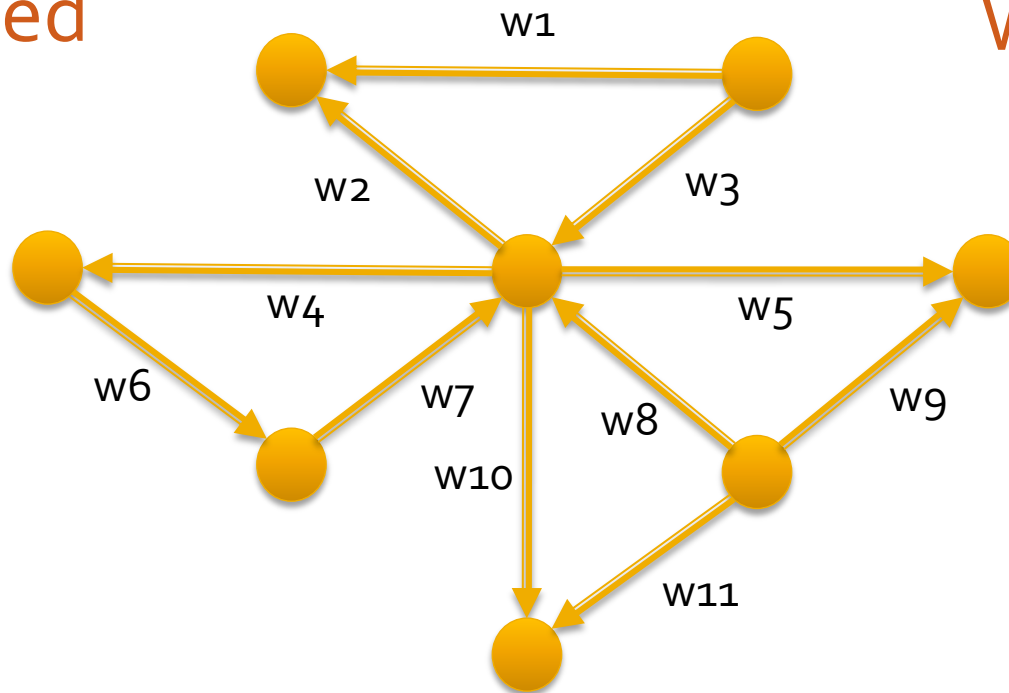
The k-Path Problem



The k-Path Problem

Directed

Weighted



Complexity

- Recursively build paths:

$$O(n^k)$$

- Can there be algorithms with run times on the form

$$f(k)n^{O(1)}?$$

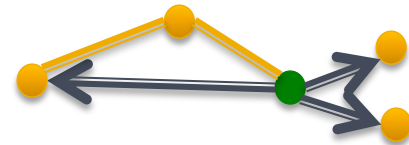
... and if so, how small can $f(k)$ be?

Intuition for FPT

Consider regular graphs of degree d :

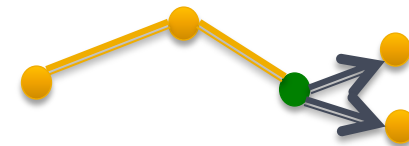
Either $d > k$:

There must be k -path.

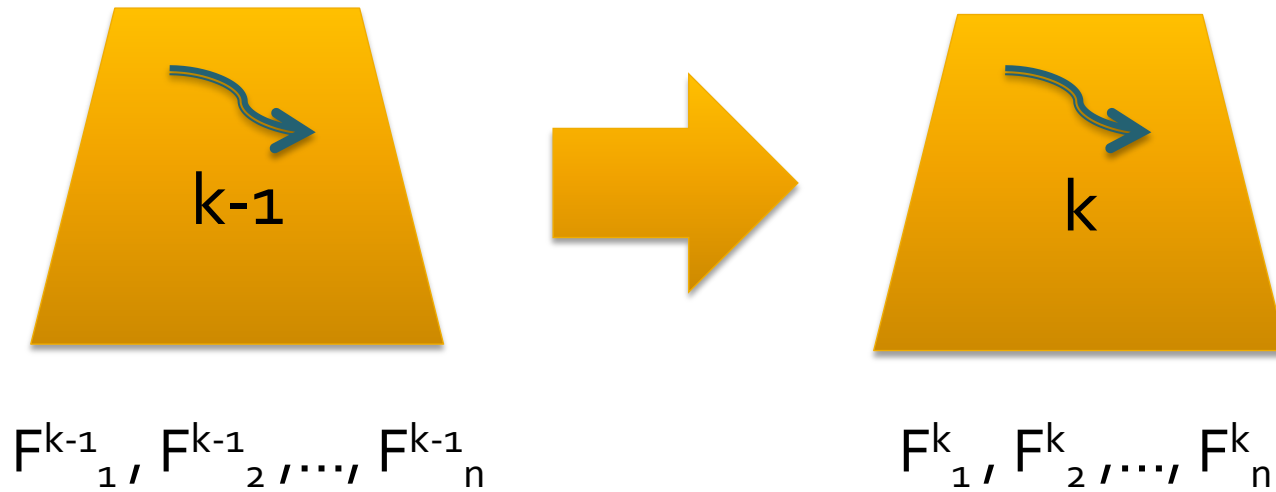


Or $d \leq k$:

We can list all potential k -paths in $nd^k \leq nk^k$ time.



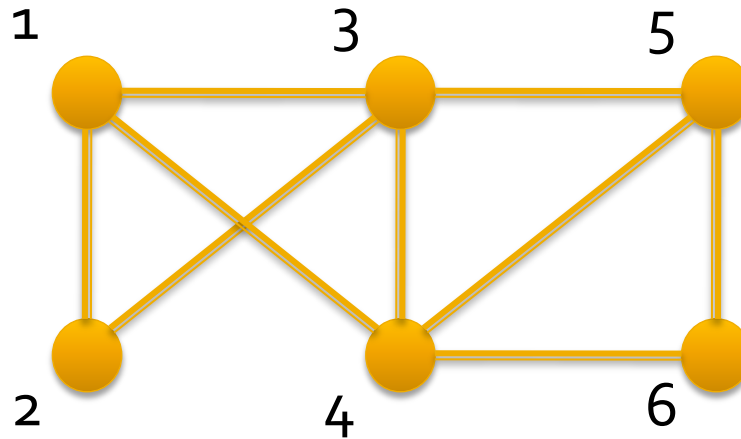
Monien 1985



$k!$

Monien 1985

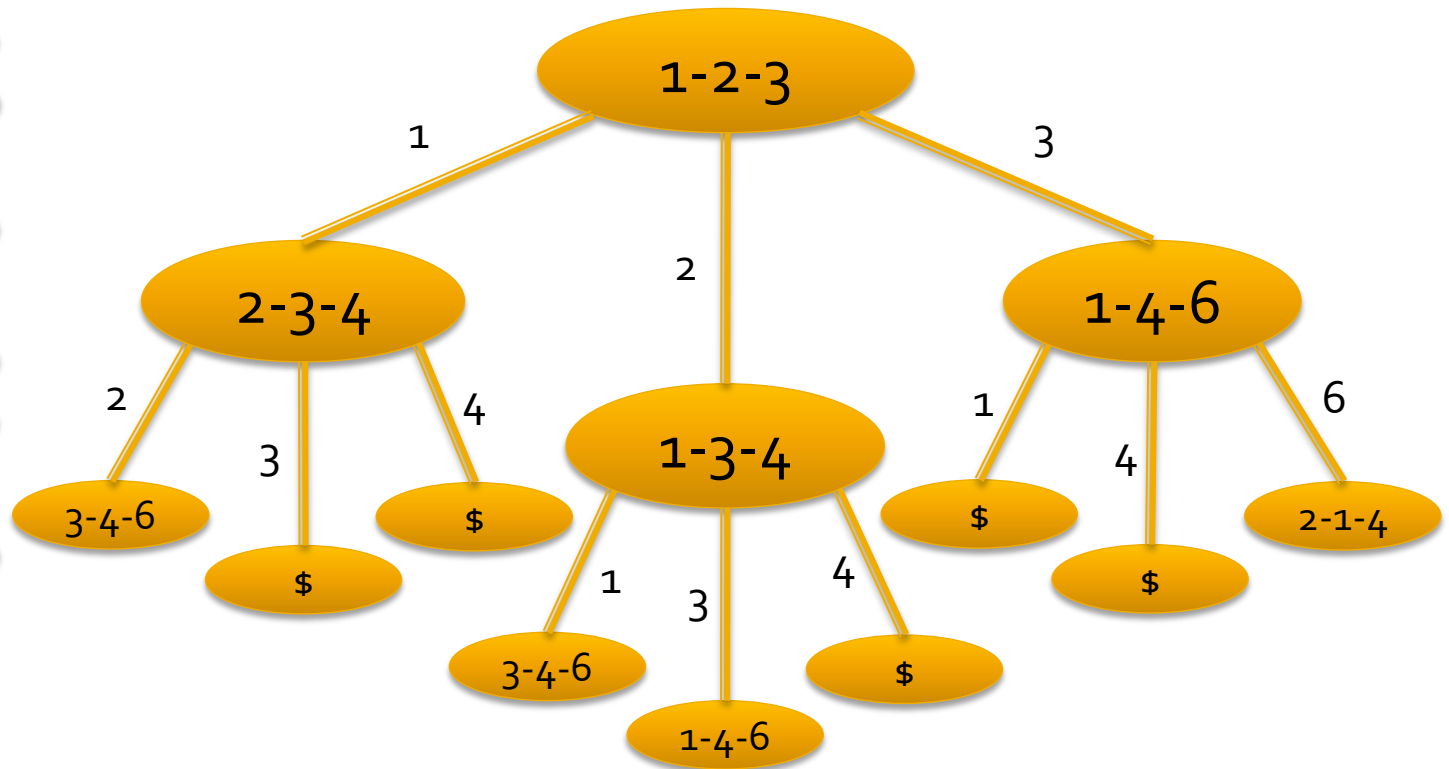
- 1-2-3-5
- 1-3-4-5
- 1-4-3-5
- 1-4-6-5
- 2-1-3-5
- 2-1-4-5
- 2-3-4-5
- 3-1-4-5
- 3-4-6-5
- 4-1-3-5



k!

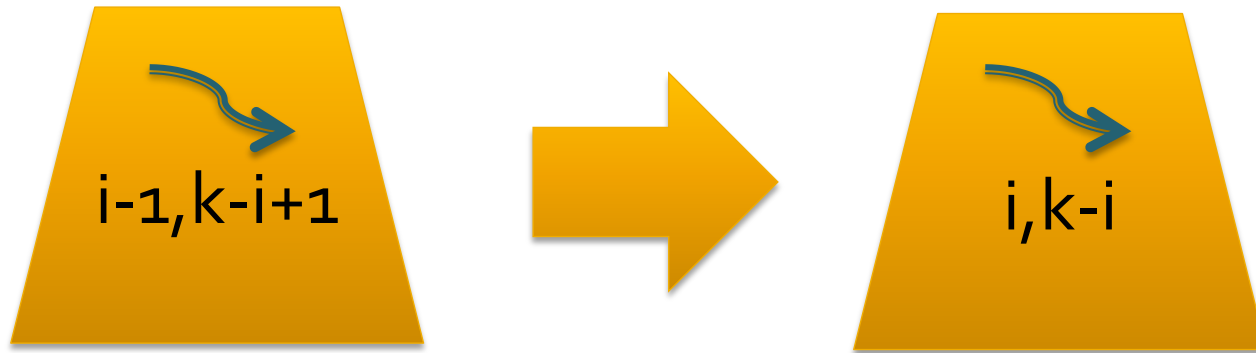
Monien 1985 Path Representatives

- 1-2-3-5
- 1-3-4-5
- 1-4-3-5
- 1-4-6-5
- 2-1-3-5
- 2-1-4-5
- 2-3-4-5
- 3-1-4-5
- 3-4-6-5
- 4-1-3-5



k!

Monien 1985

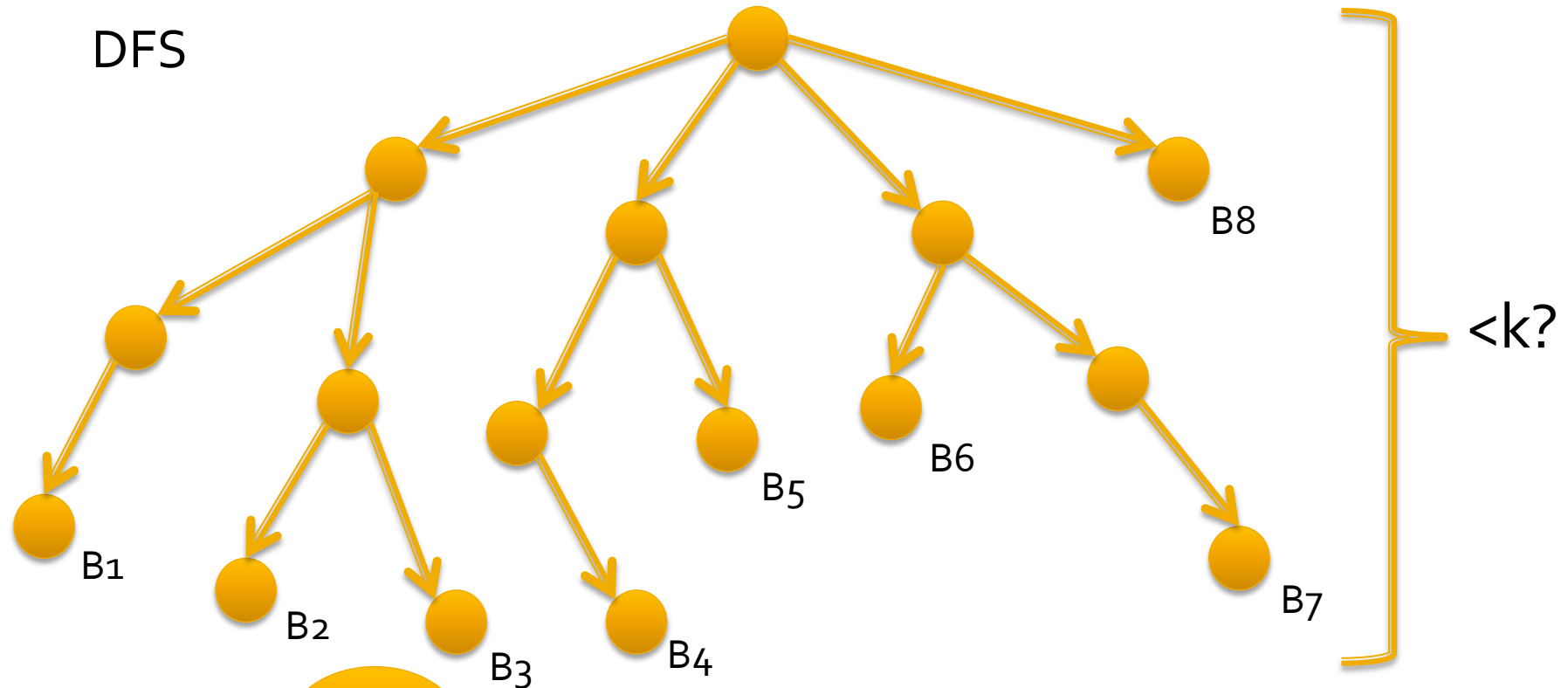


$R^{i-1, k-i+1}_1, R^{i-1, k-i+1}_2, \dots, R^{i-1, k-i-1}_n$

$R^{i, k-i}_1, R^{i, k-i}_2, \dots, R^{i, k-i}_n$

k!

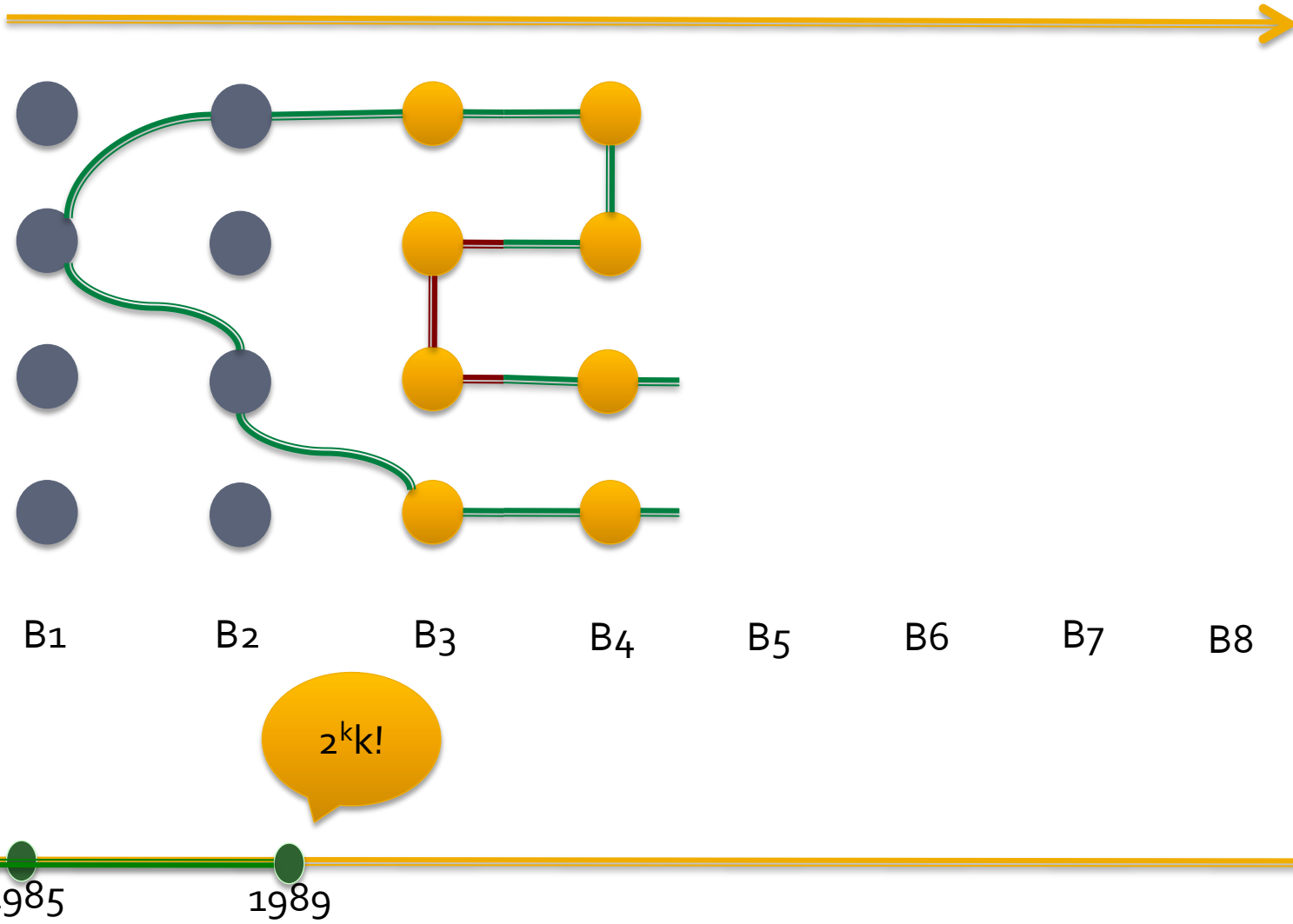
Bodlaender 1989



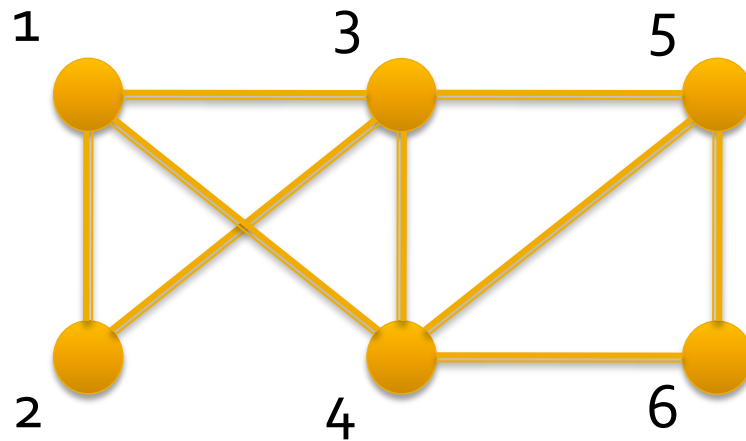
1985

1989

Bodlaender 1989



Alon, Yuster, and Zwick 1994



$\text{Prob}[\text{rainbow } k\text{-path}] \geq k!/k^k \sim e^{-k}$

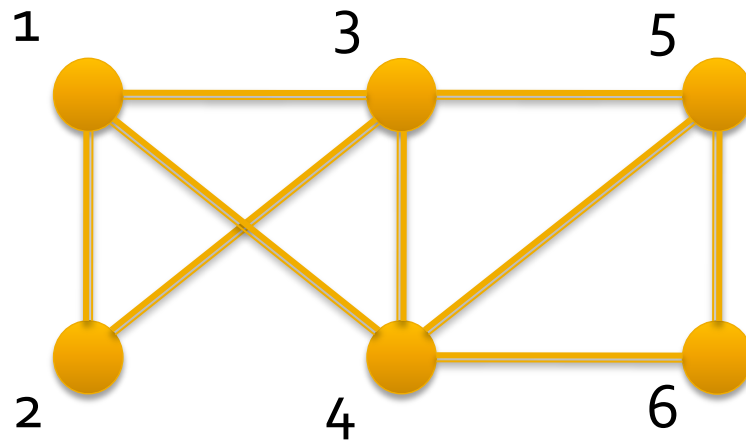
$(2e)^k$

1985

1989

1994

Kneis et al. & Chen et al. 2006



Prob[k-path split] $\geq 2^{-k}$

4^k

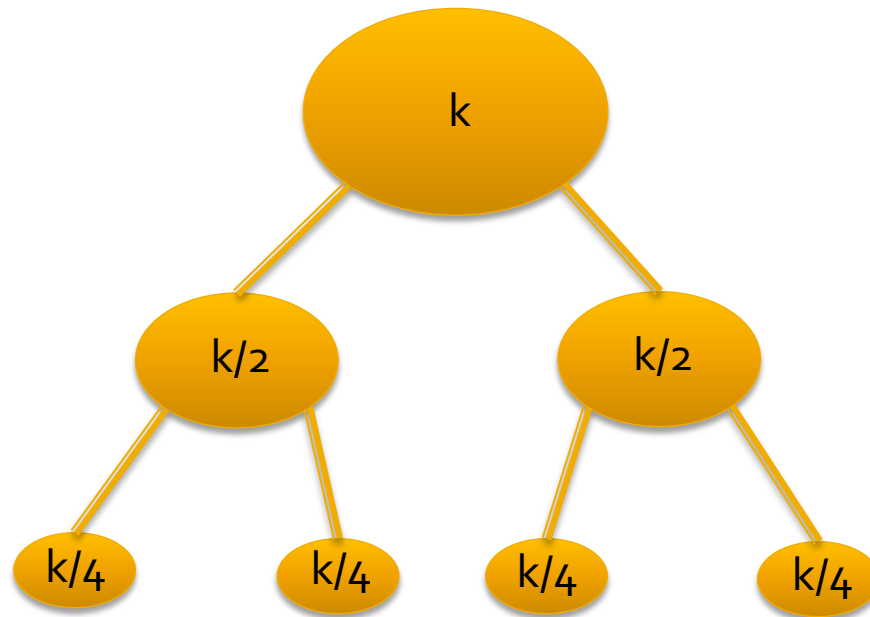
1985

1989

1994

2006

Kneis et al. & Chen et al. 2006



4^k

1985

1989

1994

2006

Koutis 2008

Combinatoric methods: Tries to construct the object explicitly piece-by-piece.

Algebraic methods: Implicitly sieves for the object by evaluating a sum.

Example: Triangle detection in a graph via fast matrix multiplication.

1985

1989

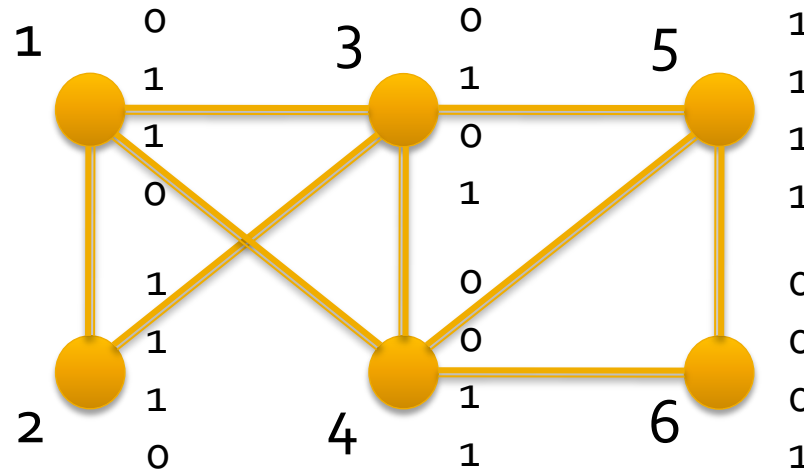
1994

2006

2008

2.83^k

Koutis 2008



Prob[k -path linear independent] $\geq 1/4$

2.83^k

1985

1989

1994

2006

2008

Koutis 2008

1	0	0	1
1	1	1	1
1	1	0	1
0	0	1	1

Prob[k-path linear independent] \geq

1 (1-1/2) (1-1/4) (1-1/8) ≥ 0.28879

2.83^k

1985

1989

1994

2006

2008

Koutis 2008

- For each vector s in $\{0,1\}^k$, look at the subgraph induced by V_s : the vertices whose vectors v_i obey $\langle v_i, s \rangle = 0$.
- Count the k -walks in $G[V_s]$, call the result w_s .
- Sum over all w_s , if sum odd report existence of k -path.

1985

1989

1994

2006

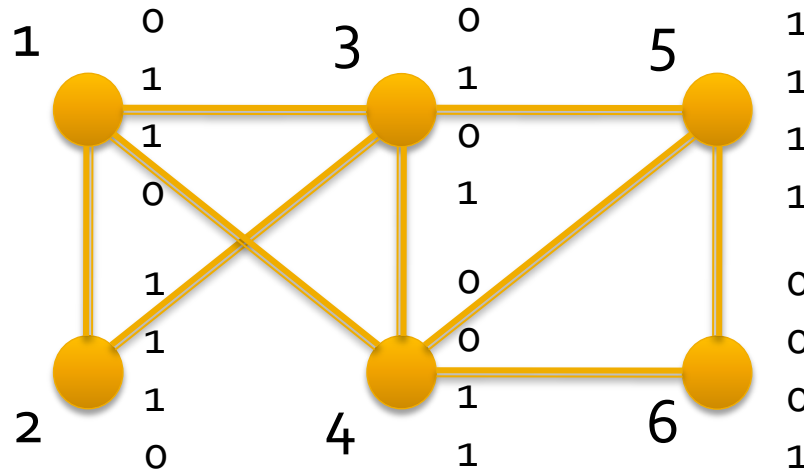
2008

2.83^k

Koutis 2008

No False Positives!

k-walk: 1,2,3,1



$$S = \begin{pmatrix} 0 & 0 \\ 0 & 1 \\ 0 & 1 \\ 0 & 1 \end{pmatrix}$$

2.83^k



Koutis 2008

Problem: Often False Negatives!

It doesn't work when all k -sized vertex sets induce a graph with an even number of Hamiltonian cycles!

1985

1989

1994

2006

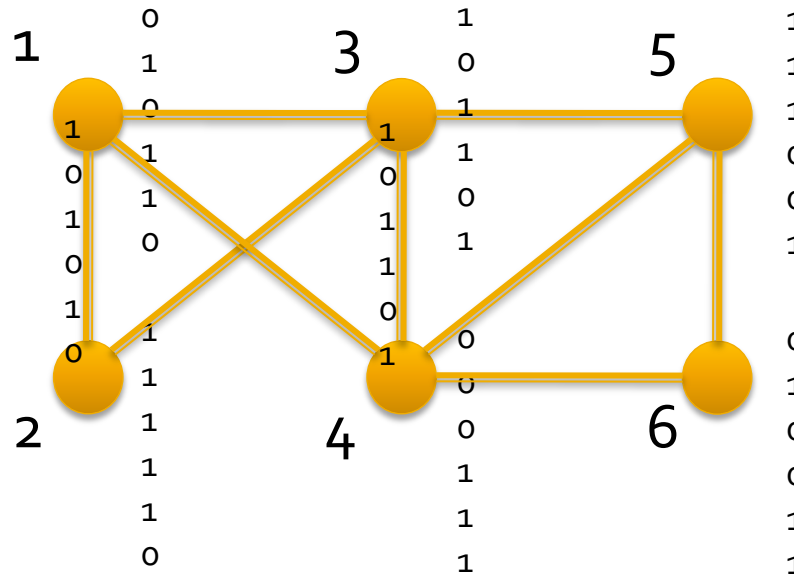
2008

2.83^k

Koutis 2008

Fix: Increase dimension and label edges as well...

k-walk: 1,2,3,4



2.83^k



Williams 2009

- Schwartz-Zippel Lemma: Let $P(x_1, x_2, \dots, x_n)$ be a multivariate non-zero polynomial of total degree d over a finite field F . For uniformly and independently sampled values r_1, r_2, \dots, r_n in F :

$$\Pr[P(r_1, r_2, \dots, r_n) = 0] \leq \frac{d}{|F|}$$

1985

1989

1994

2006

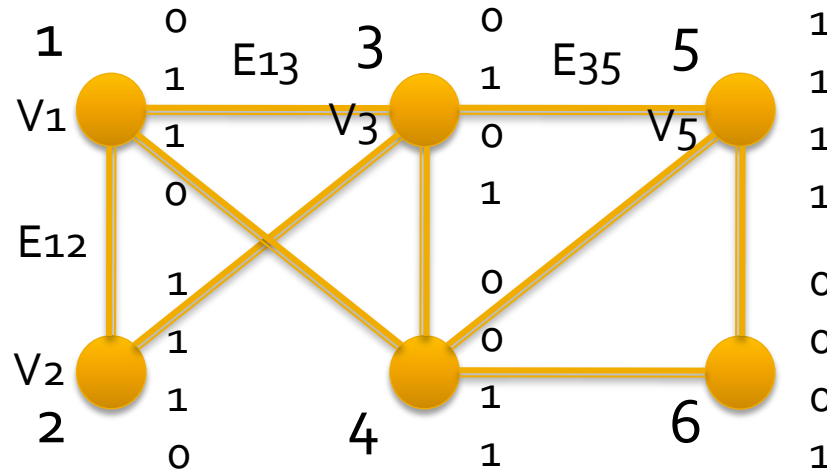
2008

2009

2^k

Williams 2009

Use larger field of characteristic two!



2^k

1985

1989

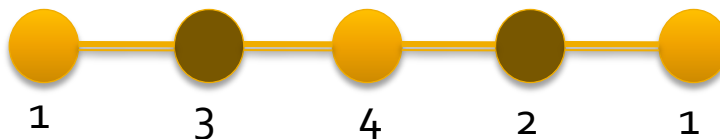
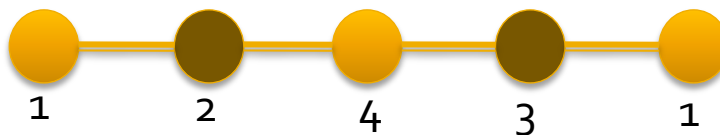
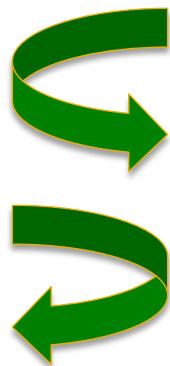
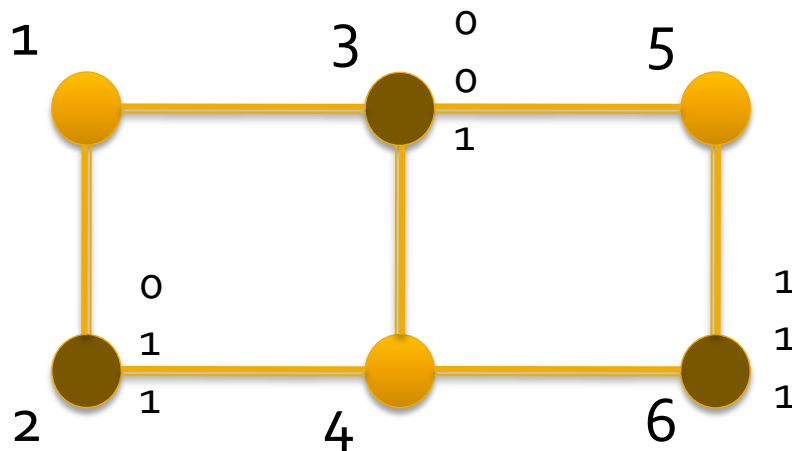
1994

2006

2008

2009

B., Husfeldt, Kaski, Koivisto 2010



$$\sum = 0$$

1.66^k



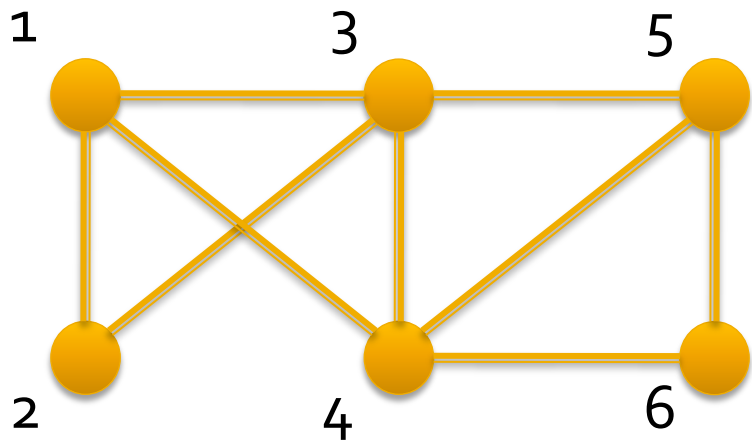
Counting k-Paths

- [Flum and Grohe 2002] Counting the number of k-paths is #W[1]-Hard.
- Approximating the number of k-paths within an arbitrarily small constant using the color-coding technique requires at least

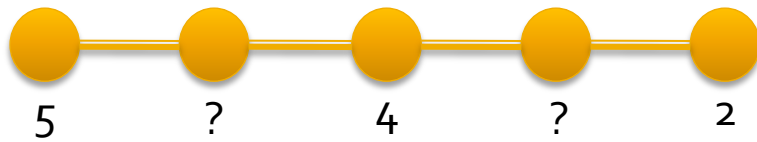
$$\Omega(n^{k/2})$$

time.

Vassilevska-Williams 2009

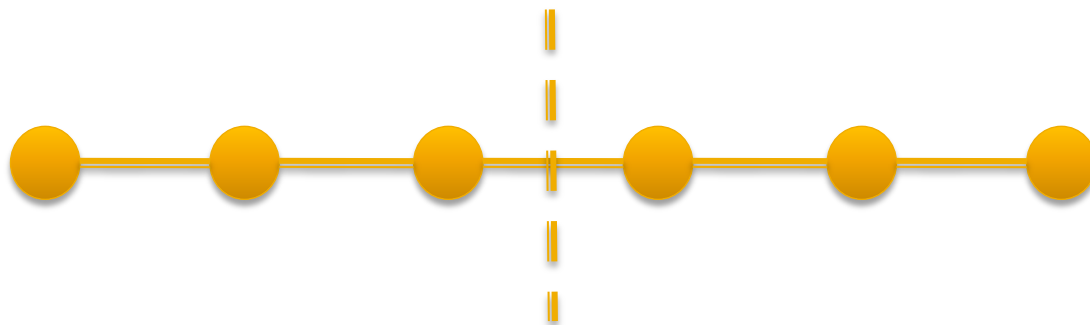


	1	3	6
5-4	0	1	1
4-2	1	1	0



$(2n)^{k/2}$

B., Husfeldt, Kaski, and Koivisto 2009



$T(S)$ = Number of $k/2$ -paths containing the vertices S .

$$\sum_{\substack{S \in 2^{[n]} \\ |S| = k/2}} (-1)^{|S|} T(S)^2$$

$$\frac{n^{k/2}}{(k/2)!}$$

Thank you for listening!