

# A QCQP Approach to Triangulation

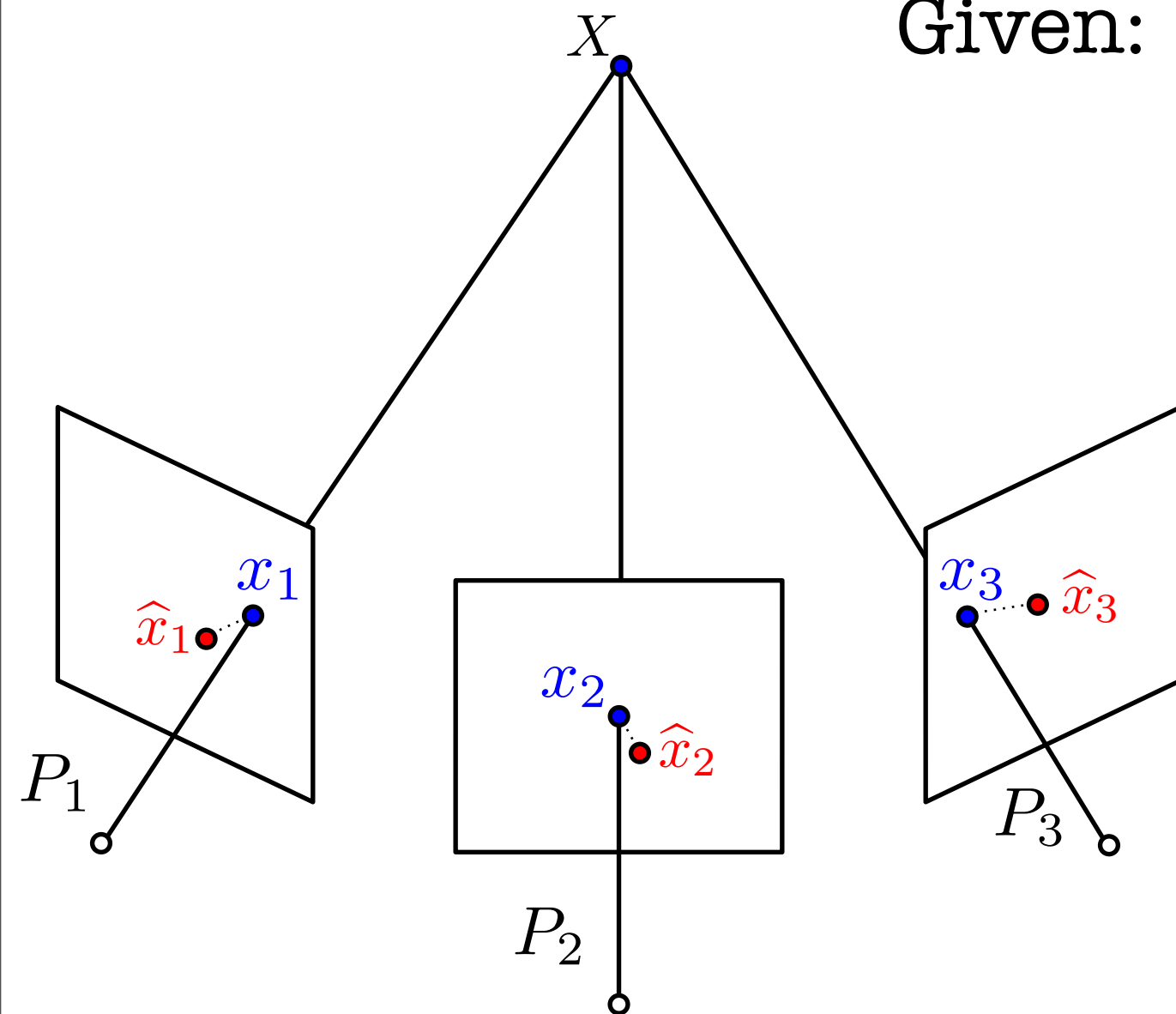
Chris Aholt<sup>1</sup>, Sameer Agarwal<sup>2</sup>, and Rekha Thomas<sup>1</sup>

<sup>1</sup> University of Washington

<sup>2</sup> Google, Inc.

# THE TRIANGULATION PROBLEM

Given:  $-n$  camera matrices  $P_i \in \mathbb{R}^{3 \times 4}$   
 $-n$  noisy observations  $\hat{x}_i \in \mathbb{R}^2$



$$P_i = \begin{bmatrix} a_i^\top \\ b_i^\top \\ c_i^\top \end{bmatrix} \quad \tilde{X} = \begin{bmatrix} X \\ 1 \end{bmatrix}$$

$$\begin{aligned} \min_{X \in \mathbb{R}^3} & \sum_{i=1}^n \|x_i - \hat{x}_i\|^2 \\ \text{s.t. } & x_i = \begin{bmatrix} a_i^\top \tilde{X} & b_i^\top \tilde{X} \\ c_i^\top \tilde{X} & c_i^\top \tilde{X} \end{bmatrix}^\top \end{aligned}$$

# PREVIOUS WORK

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## First order optimality

- Linear initialization + non-linear refinement
- [Hartley, Seo 08] - Verify global optimality of local solution
- [Hartley, Sturm 97] - Two-view triangulation
- [Stewenius, et al. 05] - Groebner basis of local solution space

## Relaxations

- [Kahl, et al. 08] - Branch and bound
- [Kahl, Henrion 07] - SDP relaxations

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Our  
Method

- SDP relaxation
- Pay attention to algebraic structure
- Polynomial in  $n$

# UNCONSTRAINED TO CONSTRAINED

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$$\begin{aligned} \min_{X \in \mathbb{R}^3} \quad & \sum_{i=1}^n \|x_i - \hat{x}_i\|^2 \\ \text{s.t.} \quad & x_i = \begin{bmatrix} \frac{a_i^\top \tilde{X}}{c_i^\top \tilde{X}} & \frac{b_i^\top \tilde{X}}{c_i^\top \tilde{X}} \end{bmatrix}^\top \end{aligned}$$



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$$V = \left\{ (x_1, \dots, x_n) : \exists X \in \mathbb{R}^3 \text{ s.t. } \forall i \ x_i = \begin{bmatrix} \frac{a_i^\top \tilde{X}}{c_i^\top \tilde{X}} & \frac{b_i^\top \tilde{X}}{c_i^\top \tilde{X}} \end{bmatrix}^\top \right\}$$

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Finding the closest point from  $\hat{x}$  to  $V$ .

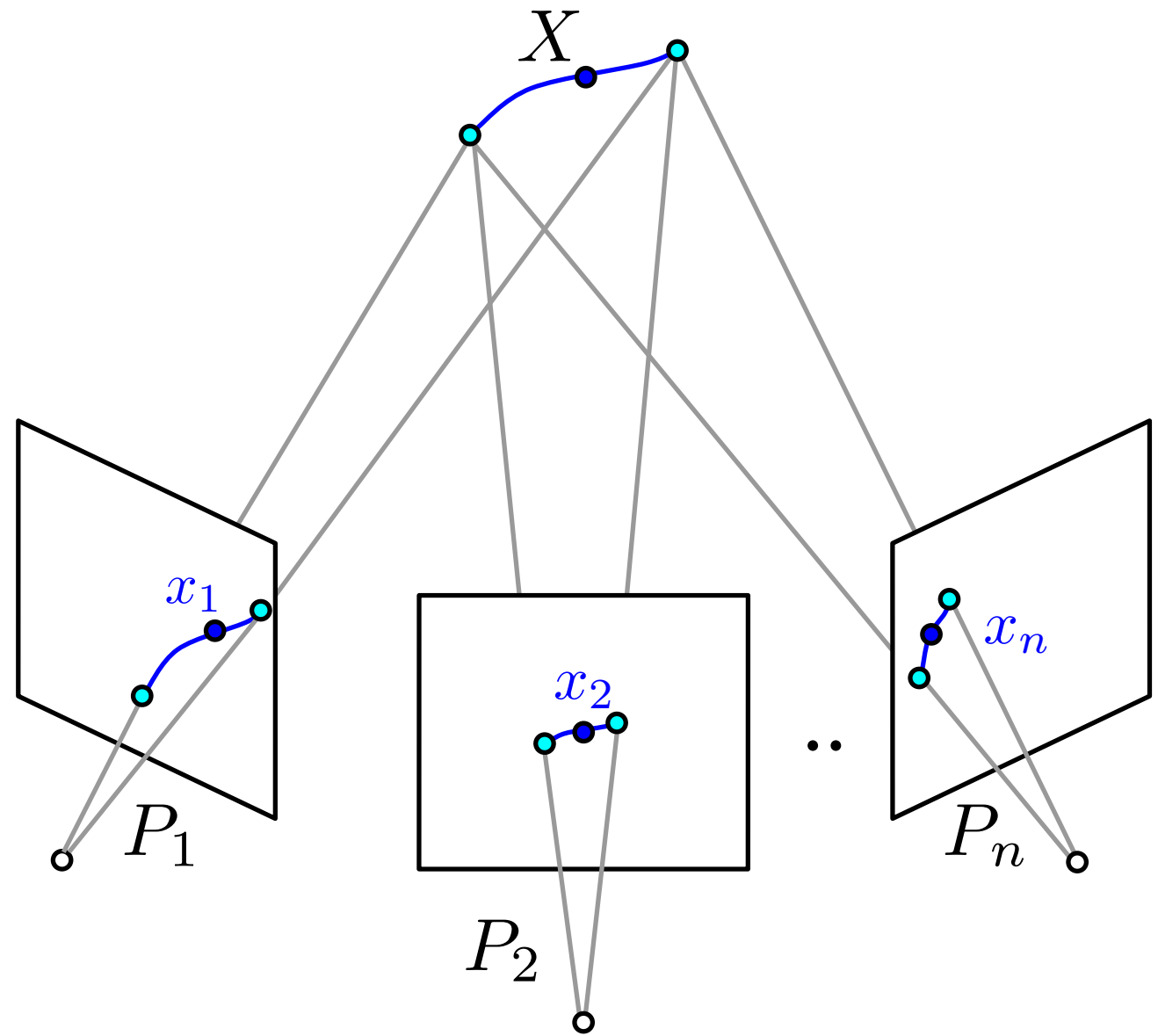
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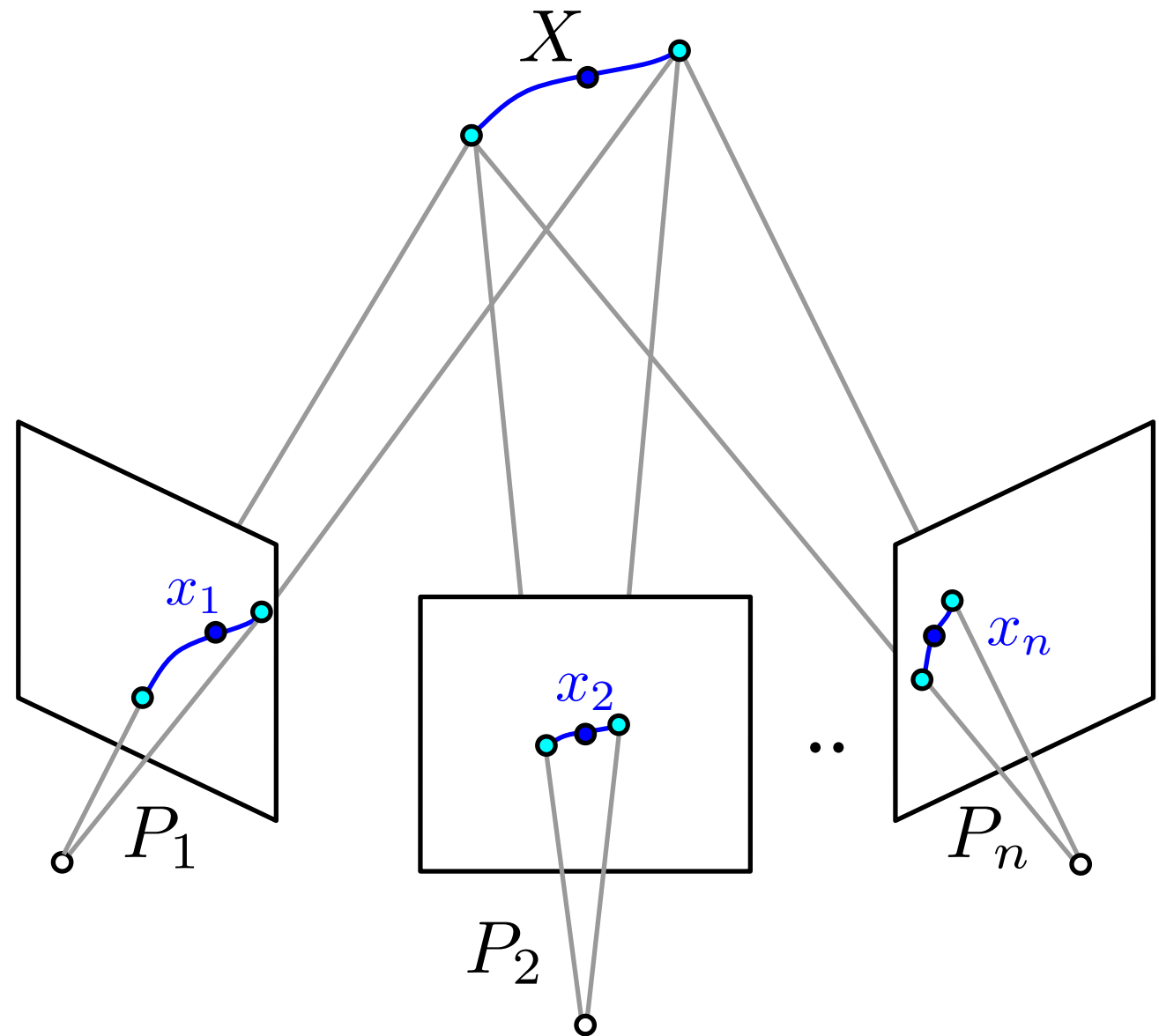
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Can we talk about  $V$  without reference to  $X$ ?

# THE CONSTRAINT SET

---

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$n > 2$  [Heyden, Åström 97]

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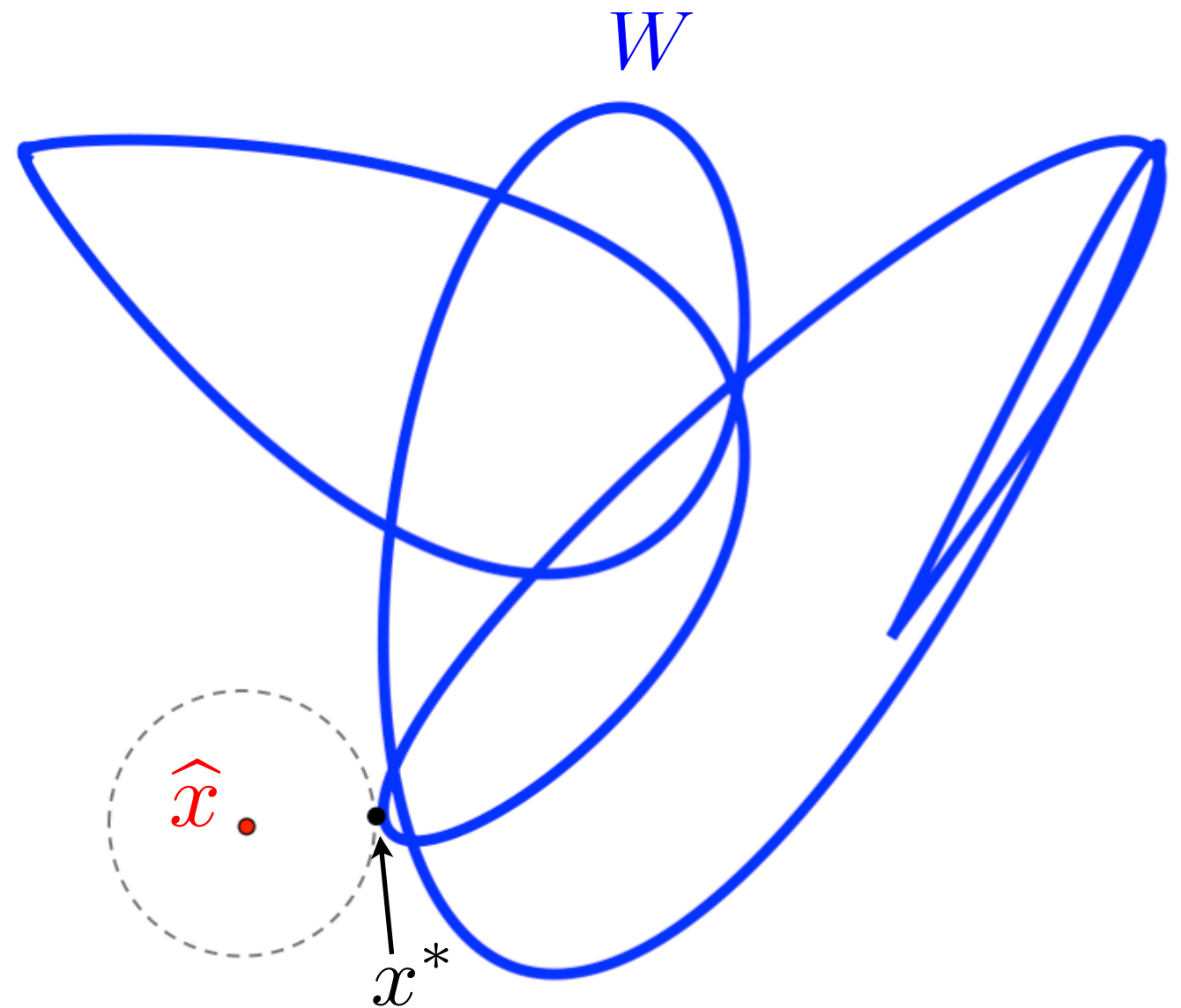
[Heyden, Åström 97]  $V = W$  when  $n = 2$  or non-coplanar cameras.

# FROM QCQP TO SDP

---

$$\begin{array}{ll} \min_x & \|x - \hat{x}\|^2 \\ \text{s.t.} & x \in W \end{array}$$

→  $x^*$



Finding the closest point from  $\hat{x}$  to  $W$ .

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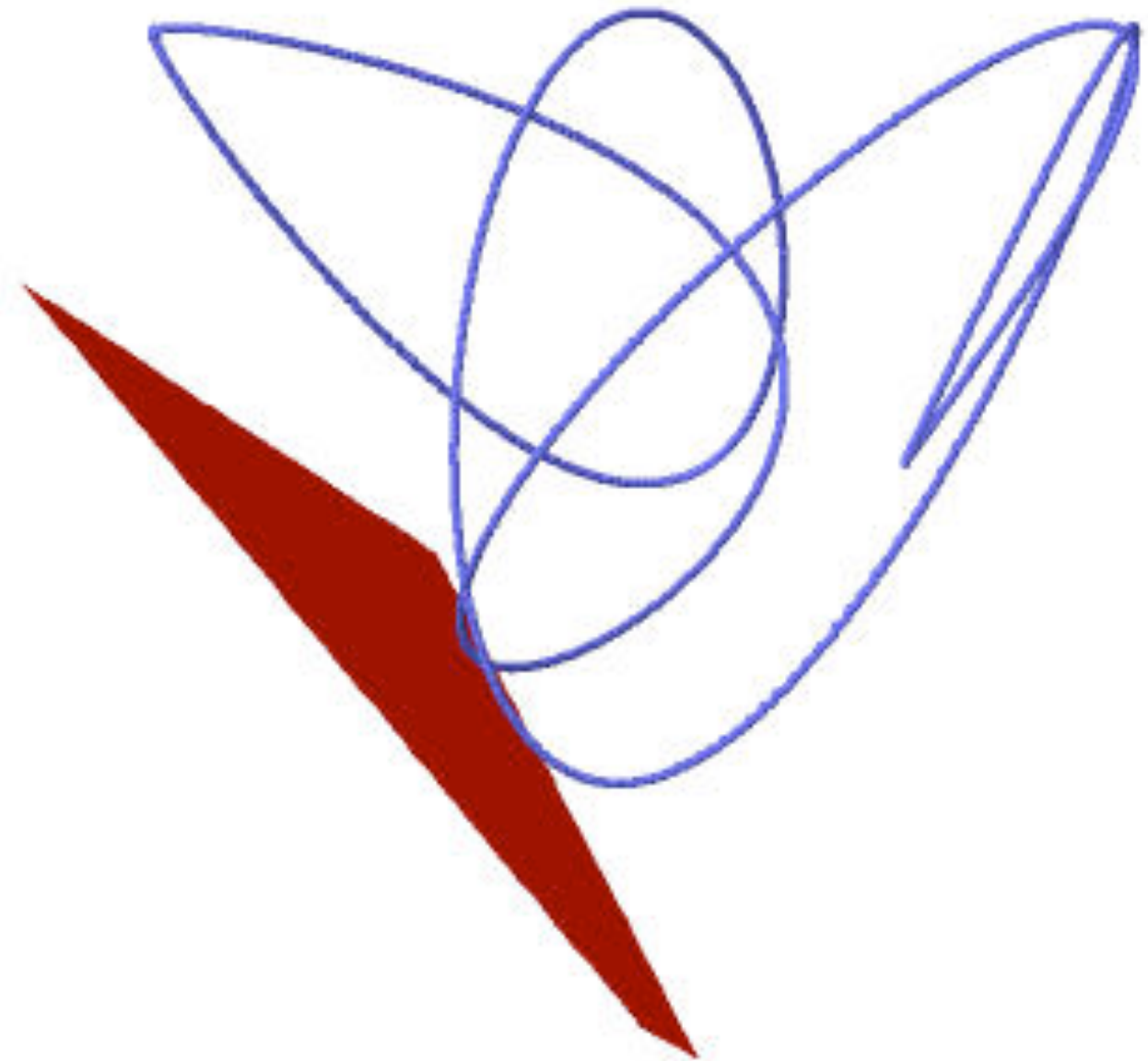
→  $x^*$

$$\min_Y \langle G, Y \rangle$$

$$\text{s.t. } \langle F_i, Y \rangle = b_i$$

$$Y \succeq 0$$

$$\text{rank}(Y) = 1$$



Rank-constrained  
semidefinite program.

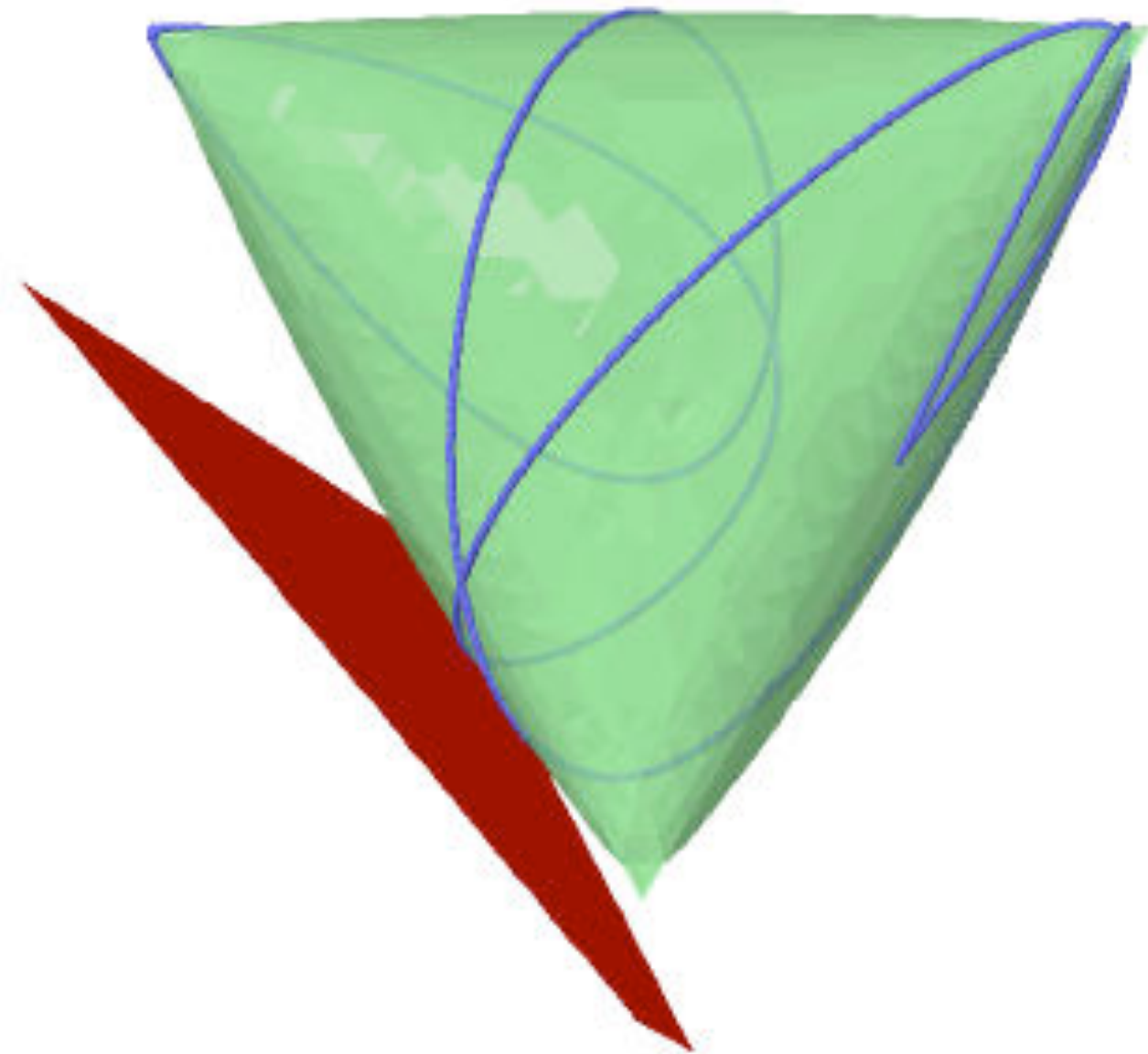
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First SDP relaxation - standard!  
Solvable in time polynomial in  $n$ .

# FROM QCQP TO SDP

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$$\text{s.t. } x \in W$$

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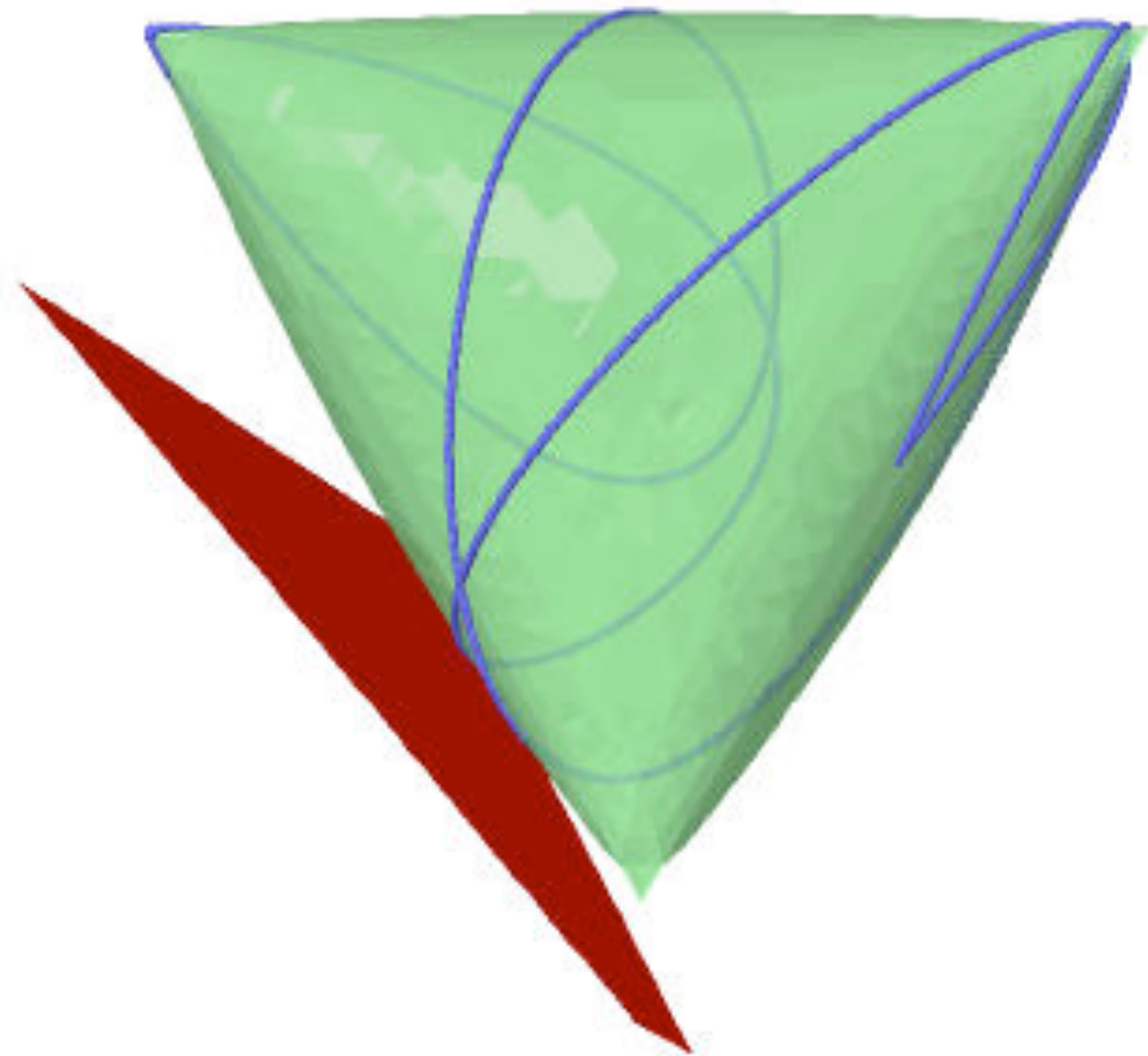
$$Y \succeq 0$$

→  $Y^*$

$$= \begin{bmatrix} \star & \\ (y^*)^\top & \star \end{bmatrix}$$

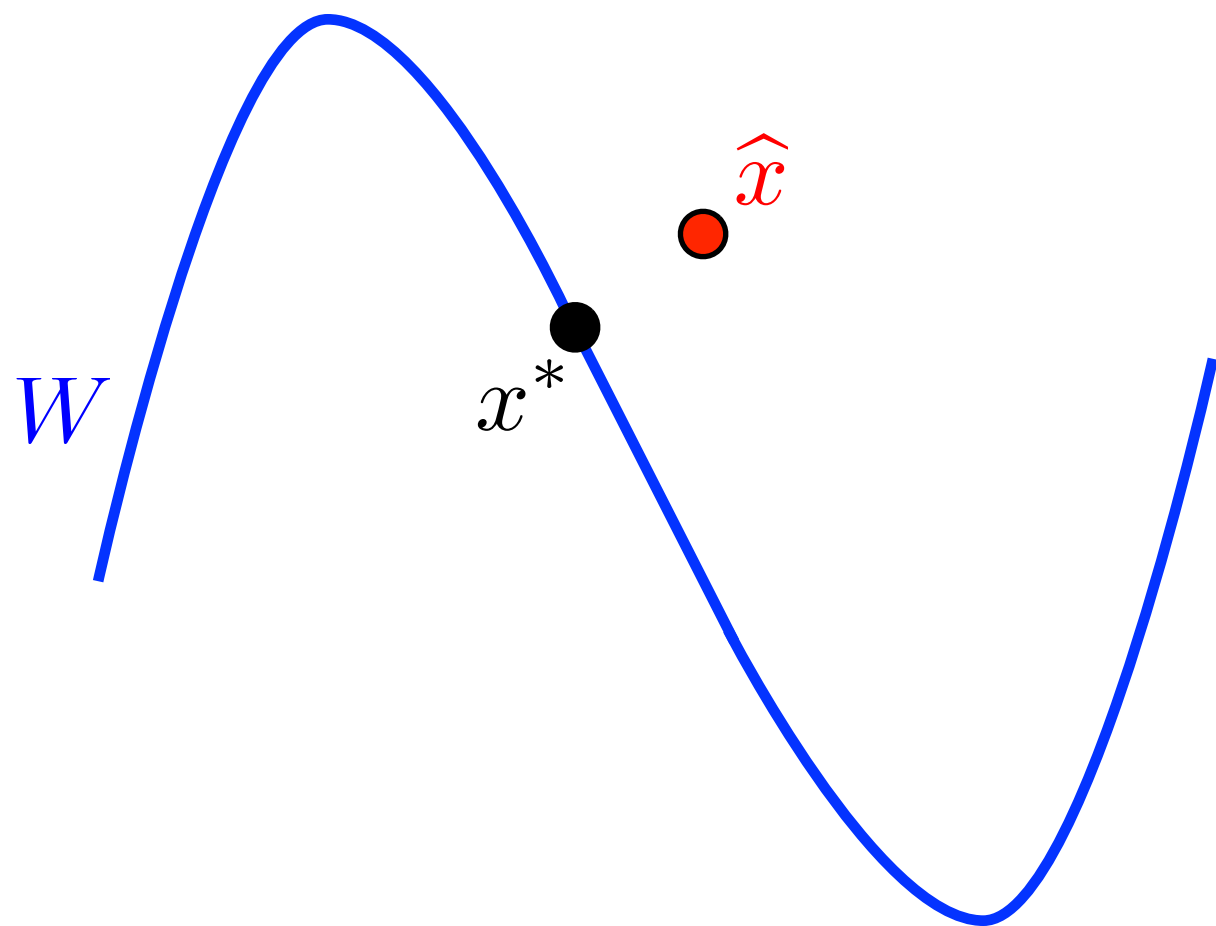
$$\begin{bmatrix} y^* \\ \star \end{bmatrix}$$

Candidate solution for QCQP



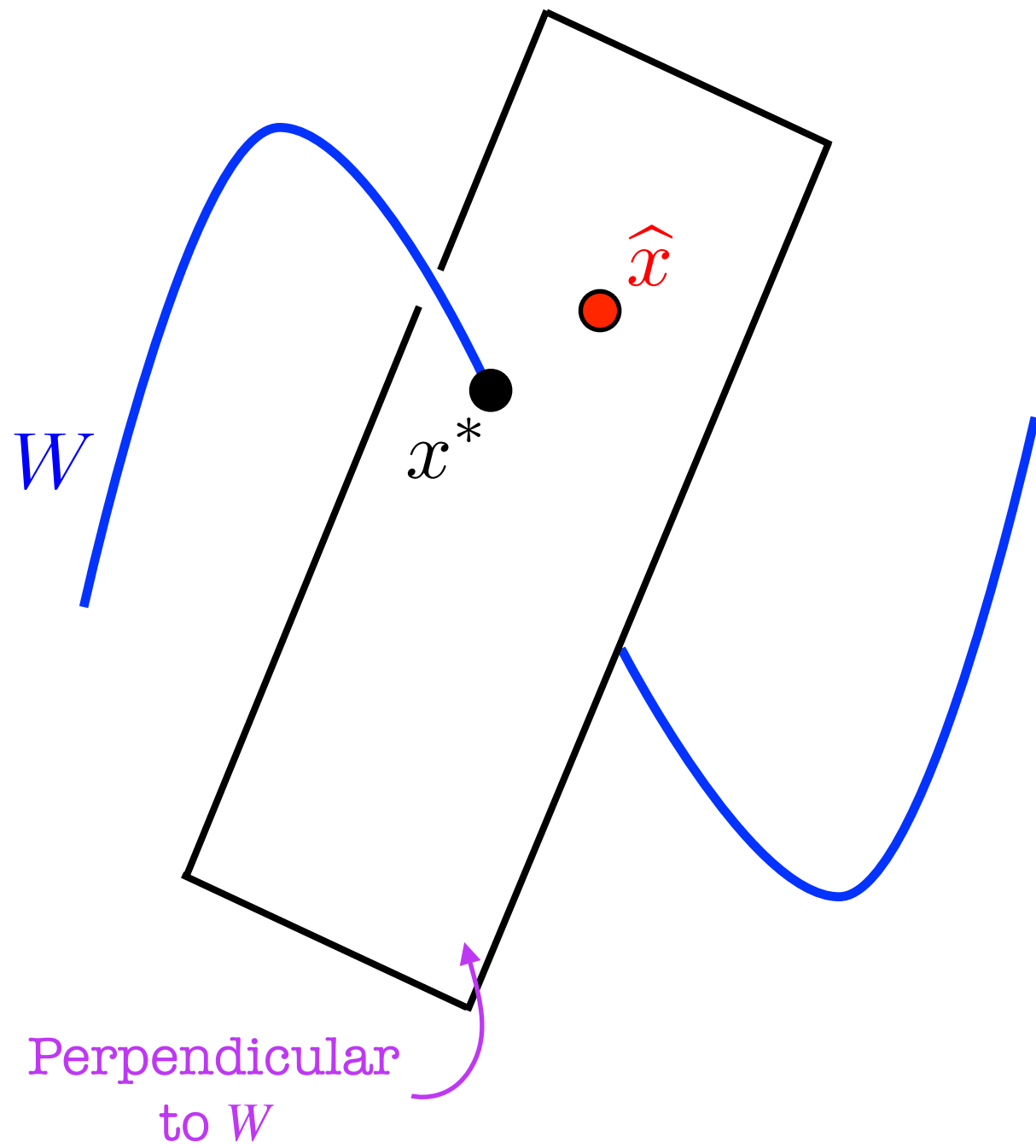
# WHEN DOES QCQP = SDP?

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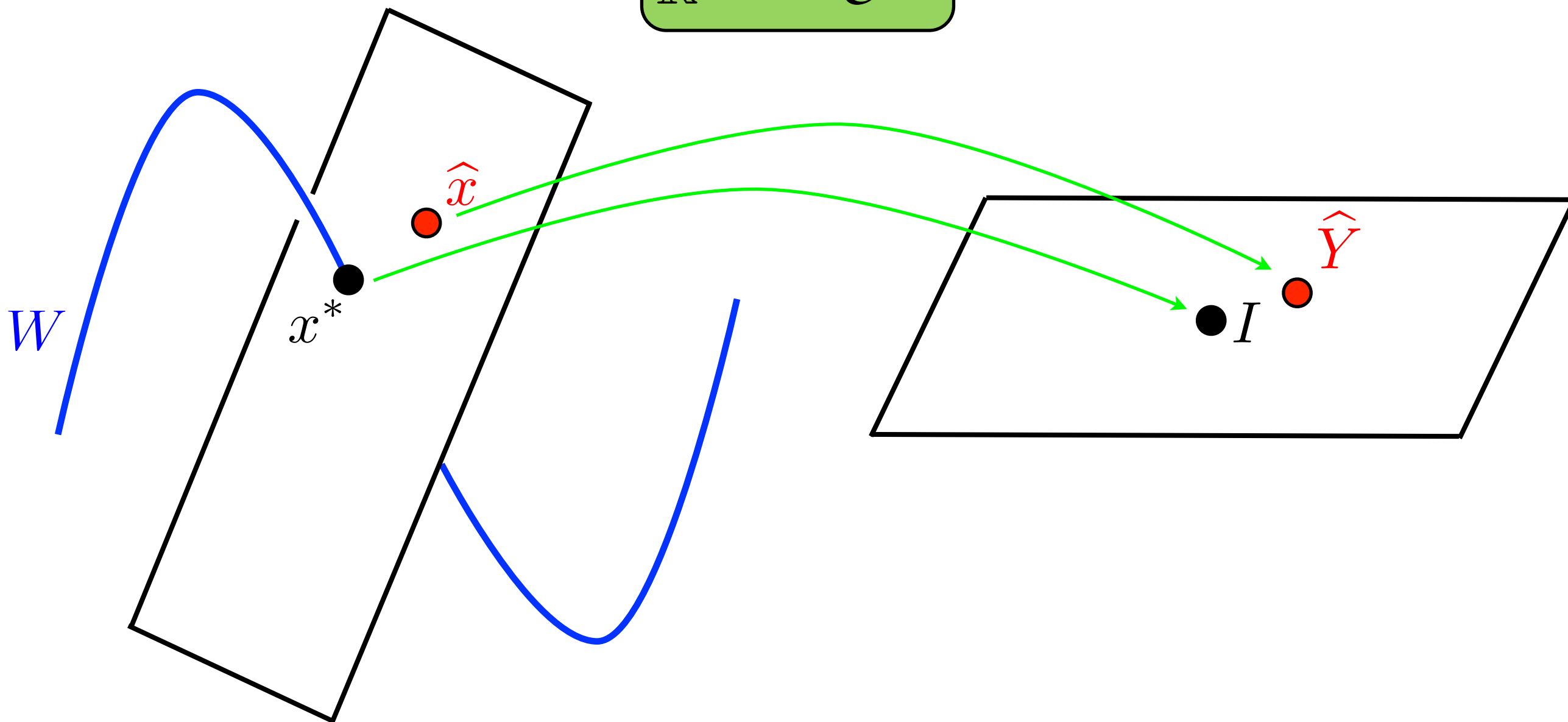
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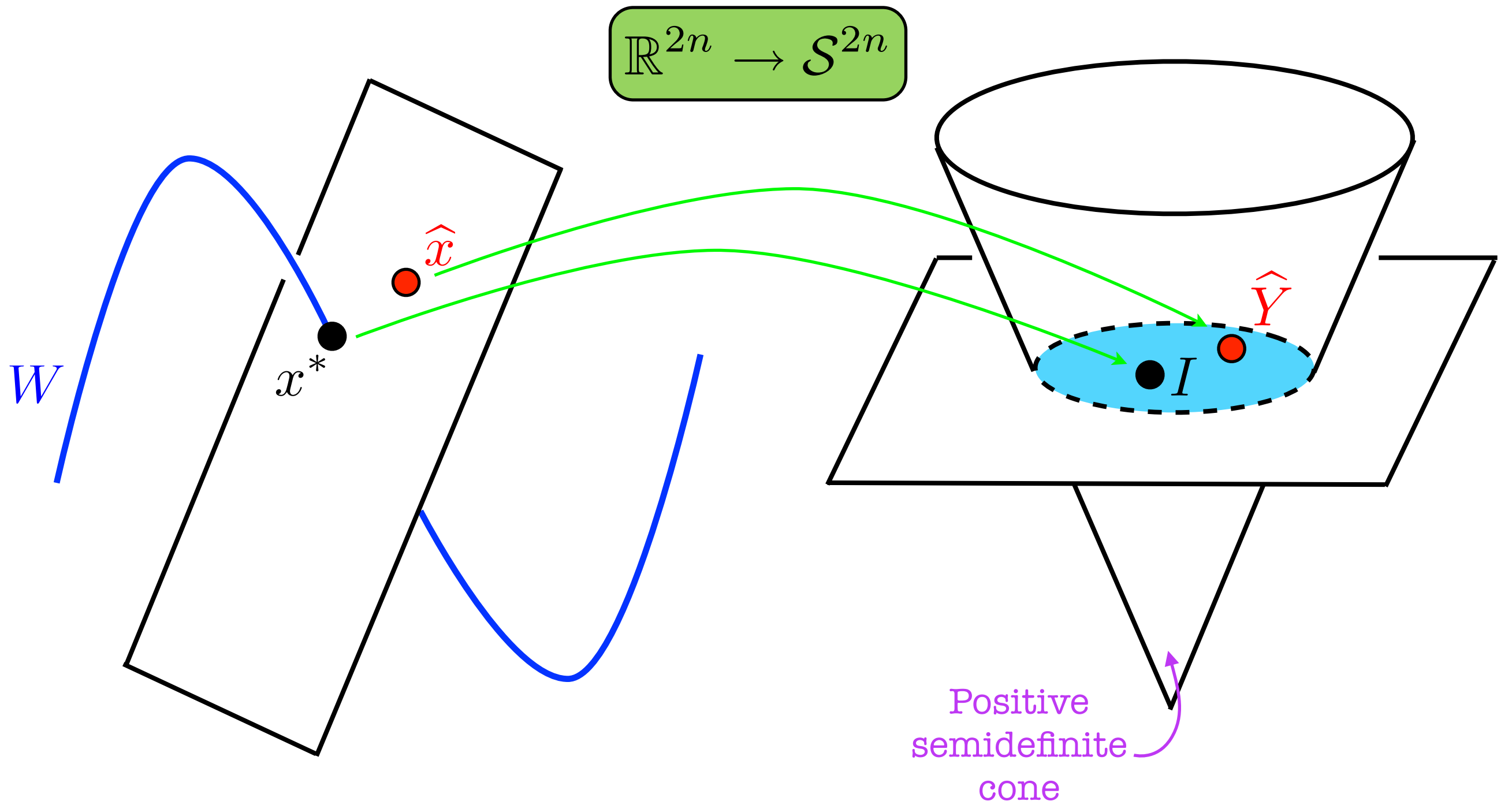
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$$\mathbb{R}^{2n} \rightarrow \mathcal{S}^{2n}$$

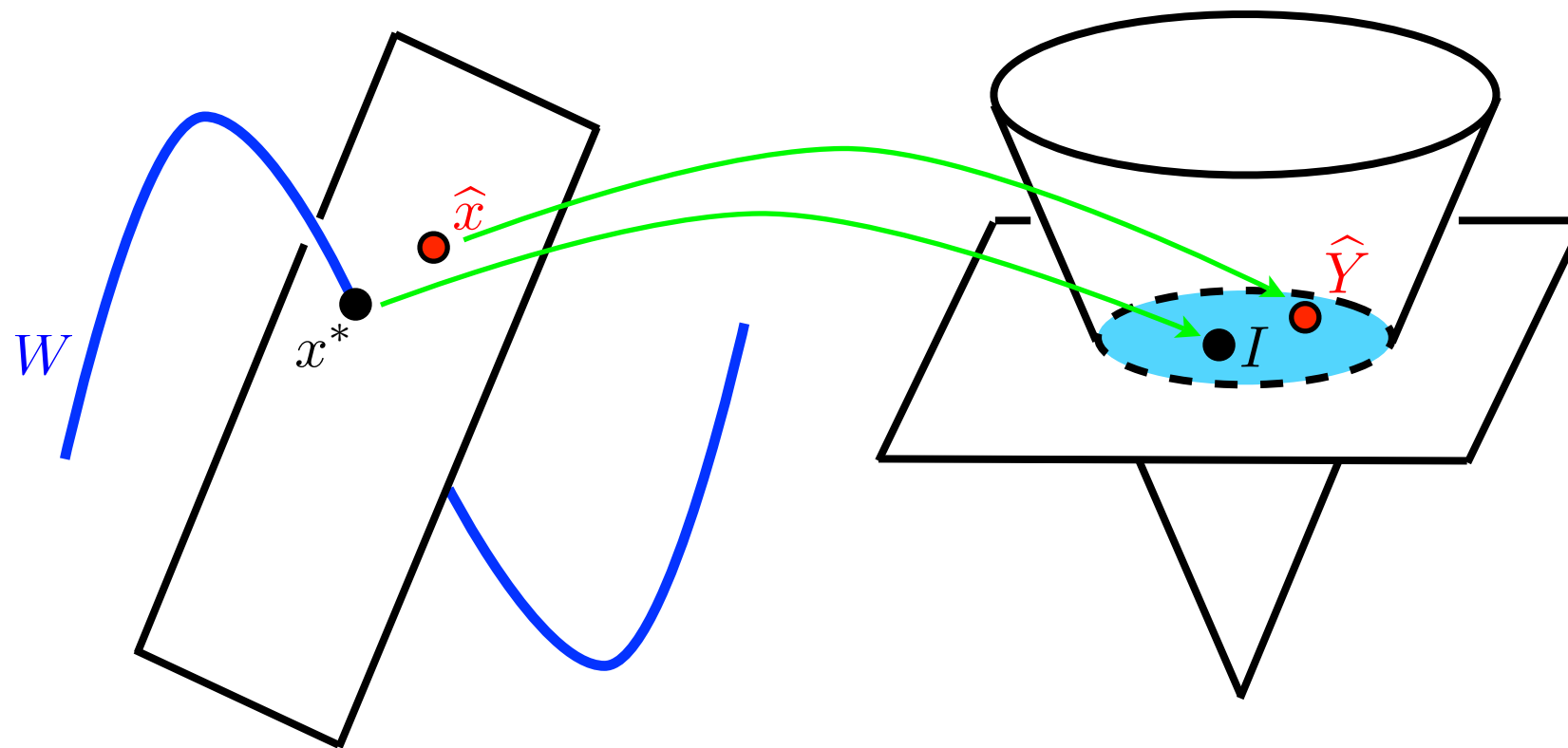




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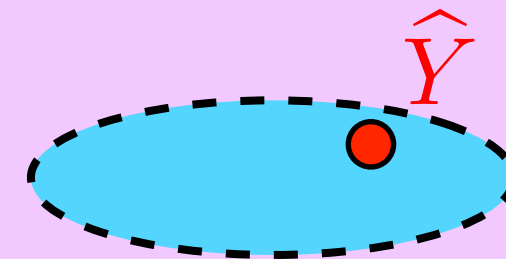
**Theorem** (A, Agarwal, Thomas)

QCQP

=

SDP

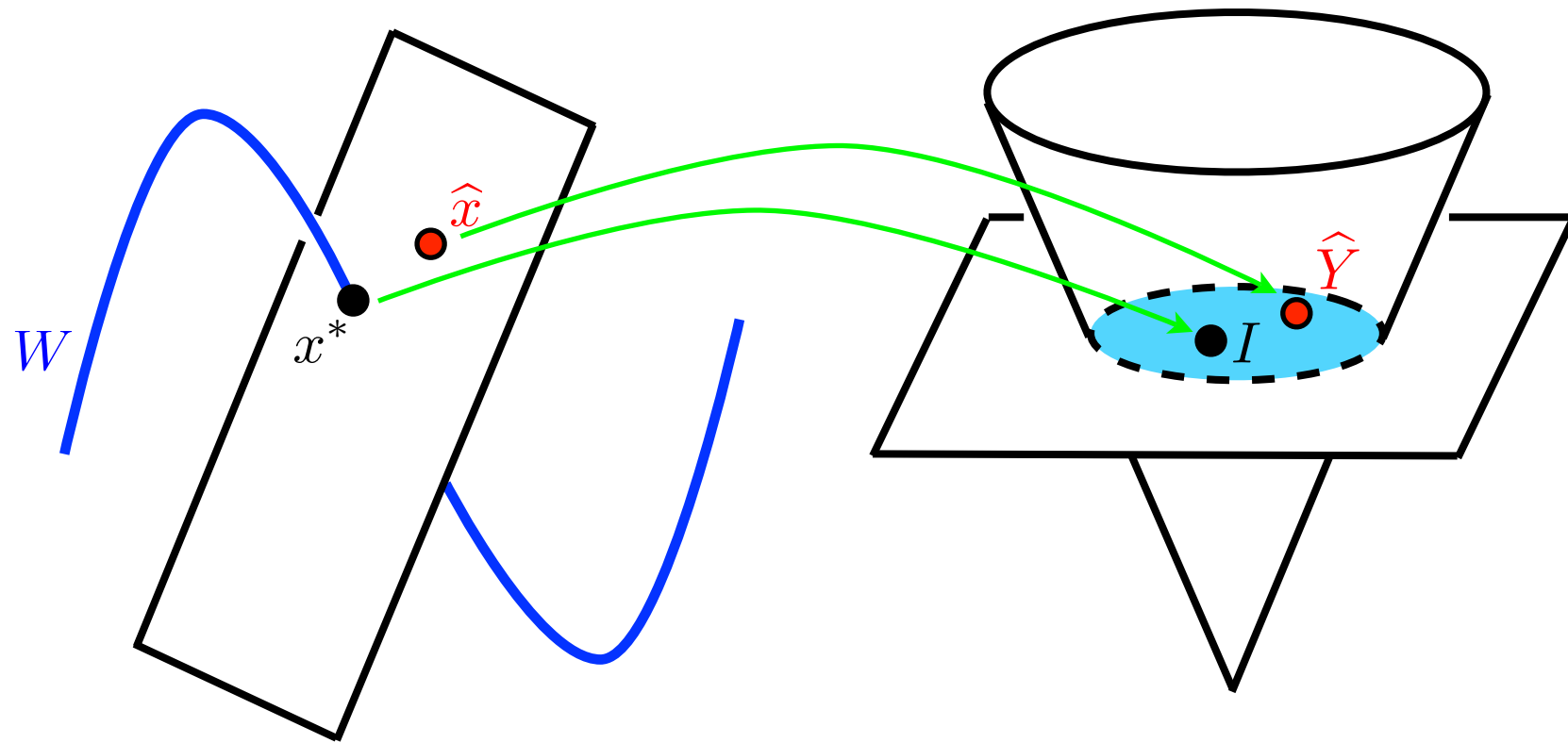
if and only if



- Can be formulated for general QCQPs with equality constraints.
- Extends previous results in special cases of QCQPs.
- Polynomial time test for optimality.

# WHEN DOES QCQP = SDP?

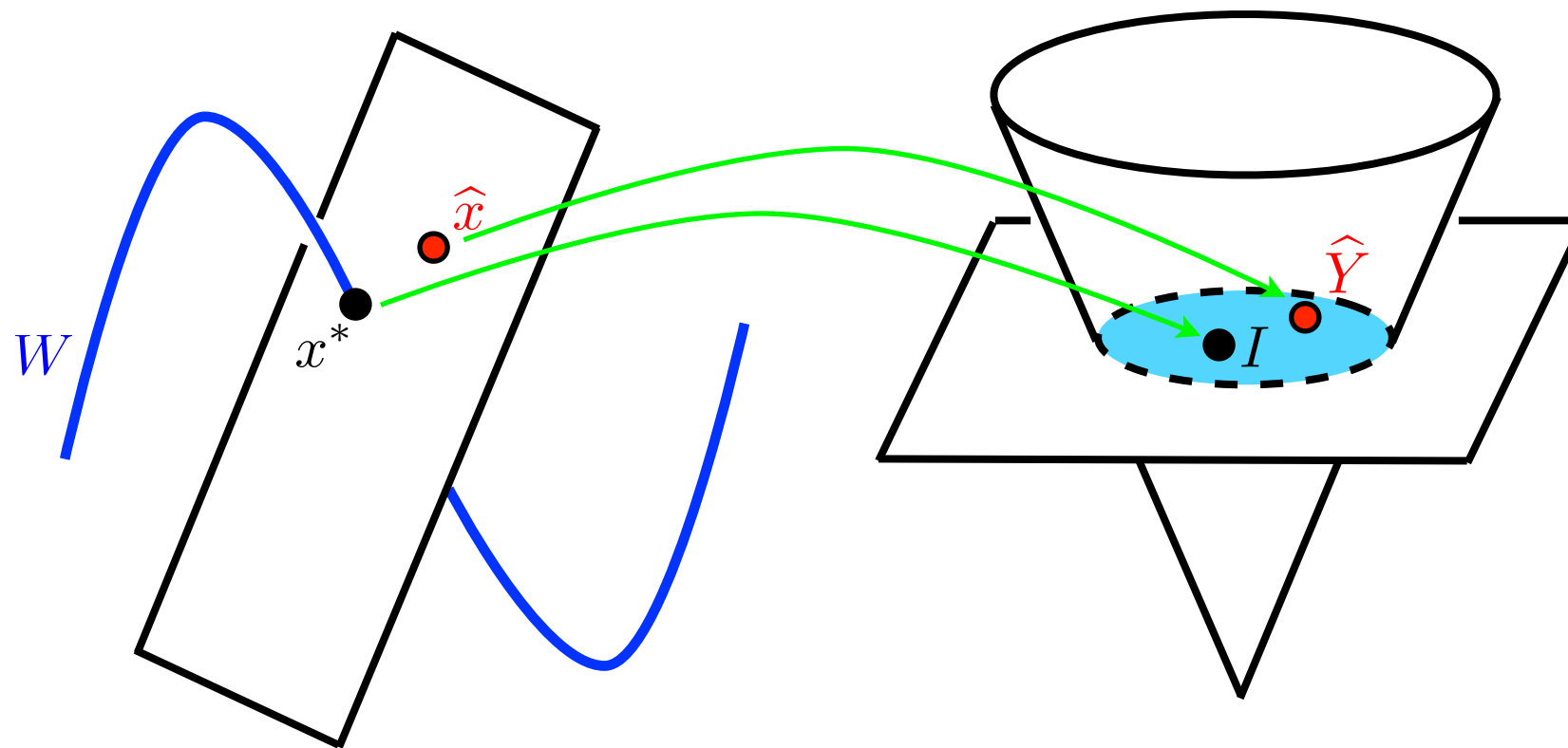
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**Theorem** (A, Agarwal, Thomas)

Always works in the case of small noise.

# WHEN DOES QCQP = SDP?



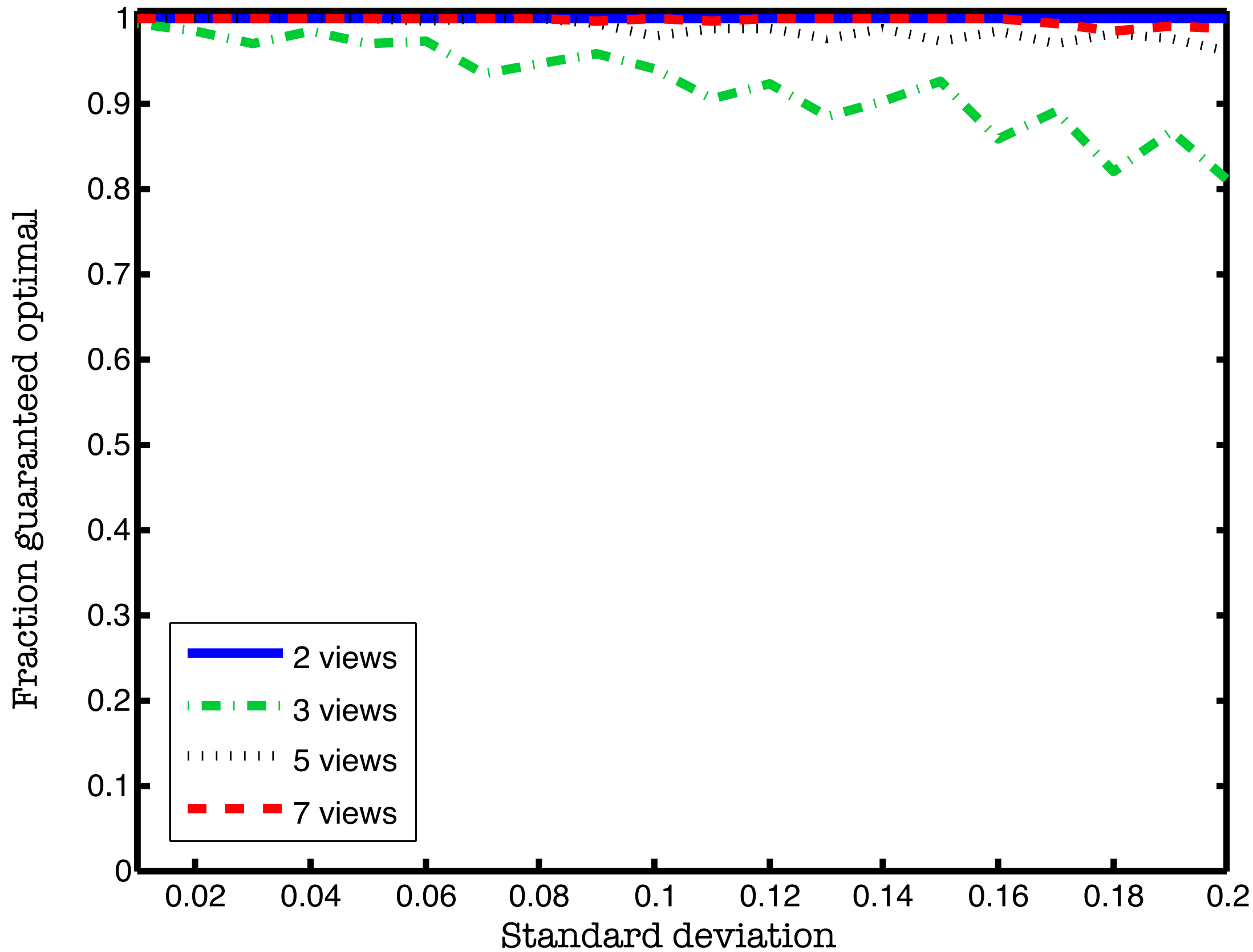
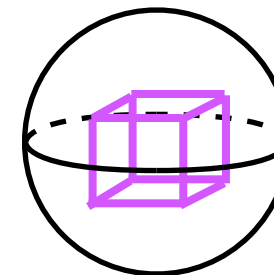
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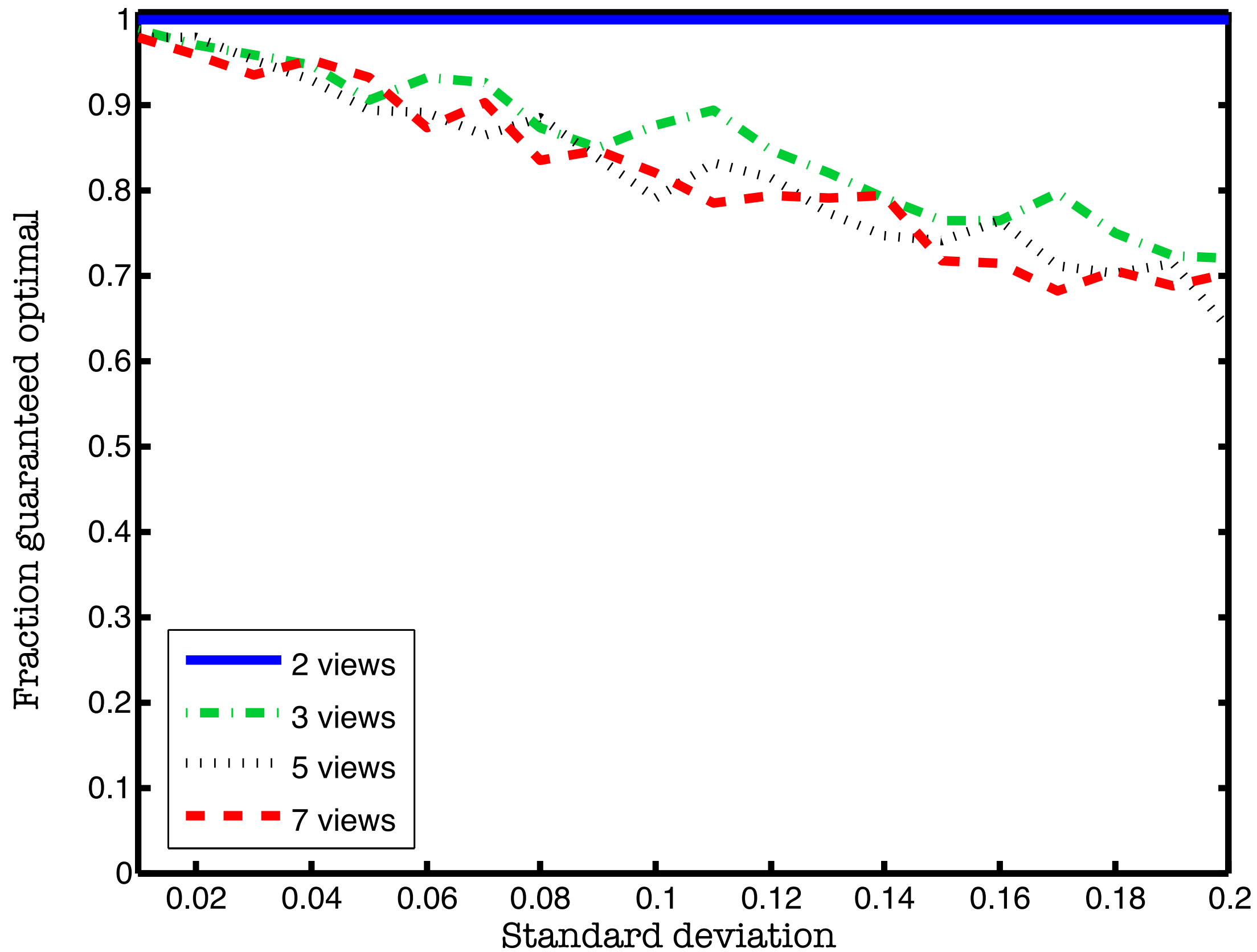
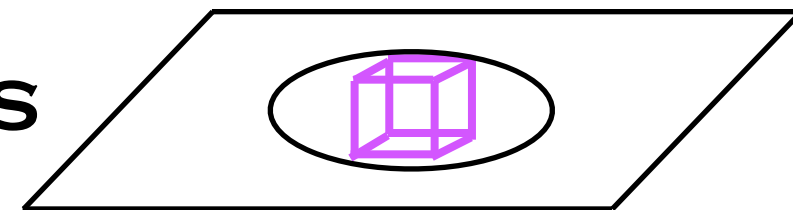
**Bonus:** Always works for two-view triangulation.

[Moré 93]

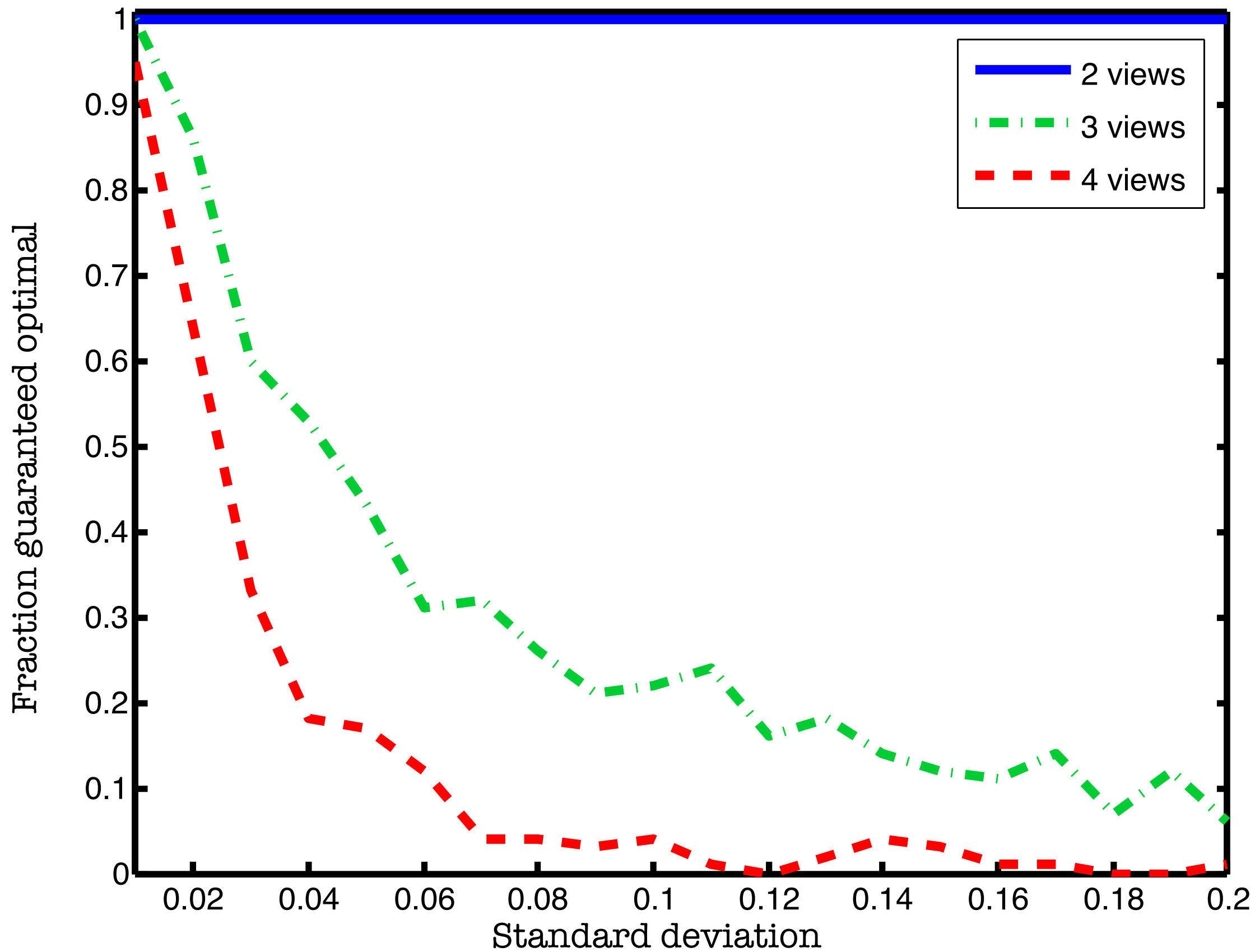
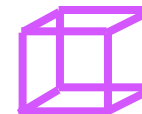
# SYNTHETIC - CAMERAS ON SPHERE



# SYNTHETIC - COPLANAR CAMERAS



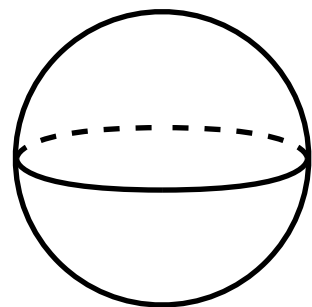
# SYNTHETIC - COLLINEAR CAMERAS



# REAL DATA

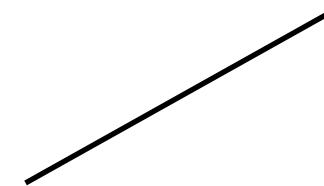
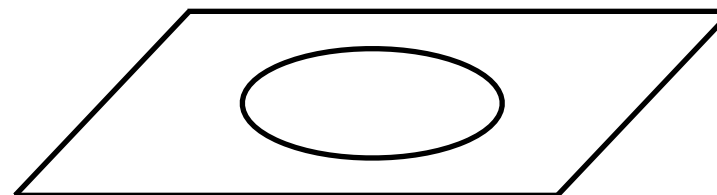
Data set	# images	# points	Optimal	Time (sec)
Model House	10	672	100%	143
Corridor	11	737	99.86%	193
Dinosaur	36	4983	100%	960
Notre Dame	48	16,288	98.4%	7200

All the camera configurations are accounted for:



Notre Dame

Model House  
Dinosaur



Corridor



# SUMMARY OF OUR CONTRIBUTIONS

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- Nearly polynomial time algorithm for triangulation.
- Geometric understanding from constraints.
- General theorems for SDP relaxations of QCQPs.