

#### RDFS Reasoning on Massively Parallel Hardware Norman Heino, Jeff Z. Pan

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#### **RDFS** reasoning

- Interpreting RDFS vocabulary can give rise to new triples
- Model-theoretic semantic conditions from RDF semantics document
- Incomplete set of rules
  - we use subset with exactly two antecedents

<pre>(5) p rdfs:subPropertyOf a (11) C rdfs:subClassOf D</pre>	<pre>q &amp; q rdfs:subPropertyOf r &amp; D rdfs:subClassOf E</pre>	$r \implies p \text{ rdfs:subPropertyOf } r$ $\implies C \text{ rdfs:subClassOf } E$
(2) <i>s p o</i>	& p rdfs:domain D	$\implies$ s rdf:type D
(3) <i>s p o</i>	& $p \text{ rdfs:range } R$	$\implies o \text{ rdf:type } R$
(7) <i>s p o</i>	& p rdfs:subPropertyOf a	$q \Longrightarrow s q o$
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• No rules with ,trivial' entailments

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- No rules with ,trivial' entailments
- Ignore non-authorative statements
- Results obtained w/o axiomatic triples

#### Special cases

(7)-1	s langProp <sub>1</sub> o	&	$langProp_1 rdfs:subPropertyOf langProp_2$	$\implies$	s langProp <sub>2</sub> o
(7)-2	s langProp o	&	langProp rdfs:subPropertyOf q	$\implies$	sqo
(7)-3	s p o	&	<pre>p rdfs:subPropertyOf langProp</pre>	$\implies$	s langProp o
(7)-4	sqo	&	$p \ rdfs:subPropertyOf q$	$\implies$	sqo

#### Special cases

(7) 1	alangDron	Q.	lang Dron ndfereuh Dronantu Of lang Dron		alangDron
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(7)	a lang Dron o	Ø.,	lang Dron ndfe Leuh Dronantur Of a		
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(7)-3	s n o	&	<i>n</i> rdfs:subPropertvOflangProp	$\implies$	s langProp o
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(7)-4	S a a	&	<i>n</i> rdfs:subPropertvOf <i>a</i>	$\implies$	S a a
(') '	590	<i>cc</i>			590

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  - MapReduce Urbani et al. (2009)
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Our platform: single host, massively parallel commodity hardware with shared memory

Challenge: develop algorithm that allows for fine-grained parallelism

#### OpenCL



- Vendor-agnostic model and API for heterogenous parallel computation
- Khronos standard
- GPU, CPU, accellerators (Cell, DSPs)
- Host program controls submission of compute kernels and data to hardware
- Compute kernels written in C99 subset



#### OpenCL memory model



1. Parsing

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- 5. Rule application on device
- 6. Copying back, storing results

#### Rule implementation



Could produce schema triples

#### Independent of (2), could be run in parallel

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#### TC-based rules

- Transitive property hierachies have a sparse adjacency matrix
  - Quadratic algorithm (Warshall) is infeasible
- Host-run serial implementation based on work by Nuutila

#### Join rules

- Join instance triple with schema subject
  - counting results
  - computing index
- Materialize result triples in two passes



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#### Duplicates by rule

		DBpedia			YAGO2 Core		
Rule	Triples	Duplicates	Ratio	Triples	Duplicates	Ratio	
(5)	0	0	_	0	19	>	
(7)	0	0	_	3,551,361	88,477	0.03	
(2)	368,832	7,630,029	21	6,450,781	13,453,038	2.1	
(3)	568,715	4,939,870	8.7	409,193	1,511,512	3.7	
(11)	259	610	2	3,398,943	366,764	0.1	
(9)	0	8,329,278	>	6,685,946	3,173,957	0.5	
(11+9)	259	10,398,328	42,162	35,061,599	57,969,000	1.7	
all	1,650,607	23,775,152	14	45,766,218	89,370,361	2.0	

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#### Global duplicate prevention

- Hash table in global device memory
- Each thread materializes triple only if not in hash table
- Table is static (not updated)

#### Local deduplication

- 1. Sort triples in local memory
- 2. Count adjacent duplicates
- 3. Caclulate new duplicate-free index
- 4. Zero out duplicates, rearrange unique values

data vector

data vector

data vector

data vector



# Local dedup example data vector index vectors



0 I 0 0 0 I 0 I 0 I I I I I *I I* 2 3 *K* 

#### Local dedup example index vectors data vector 4 5 d $\mathbf{0}$ k $r_i = \left\{egin{array}{cc} d_{i+k_i} & i+k_i < |d| \ 0 & ext{else} \end{array} ight.$

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#### Experiments

- Exp1: study scalability on different levels of hardware parallelism
  - performed on 4 CPU server with 32 CUs
- Exp2: study efficacy of duplicate removal strategies
  - performed on 20 CU GPU device
- Exp3: compare to previous work

#### Datasets

- DBedia ontology, infobox types, infobox triples:
  - ~26M triples
  - ~1.7M new closure triples
- YAGO2Core
  - ~36M triples
  - ~46M new closure triples
- Both scaled to 1/2, ..., 1/16th of instance triples plus all schema triples

#### Results – Exp1



#### Results – Exp2

Dataset	Strategy	Kernel time (ms)	Closure time (ms)	Duplicates	Speedup
DBpedia	None	28.444	6,884.15	23,775,152	
	L	120.915	6,083.76	12,165,520	13.2 %
	G	52.305	6,635.60	1,511,758	3.7 %
	L+G	117.400	6,557.94	1,057,470	5 %
YAGO2/8	None	25.565	21,625.19	31,552,221	
	L	187.169	19,554.09	2,399,898	10.6 %
	G	53.948	21,622.31	29,357,936	0%
	L+G	215.947	19,807.66	1,786,753	9.2 %

#### Results – Exp3

	Input triples	Output triples	Damásio (ms)	Our system (ms)	Speedup
T2	366,490	3,617,532	23,619.90	9,038.89	$2.6 \times$
Τ6	1,942,887	4,947,407	18,602.43	1,964.49	9.5×

T6 also performed on M/R cluster in > 3 min

#### CPU – GPU comparison

Device	Kernel execution (ms)	Total (ms)
Core i7 3770 (CPU)	647.311	5509.92
Radeon HD 7870 (GPU)	114.683	5881.54

#### Conclusions/future work

- RDFS reasoning is can be done *massively* parallel
- Shared memory can be used for efficient parallel duplicate reduction
- low ALU:fetch ratio is unvaforable for GPU devices
- Data compression
- Multiple devices
- Reasoning on the FPU
- Complete implementation w.r.t. RDF semantics

#### Complexity of RDFS reasoning

- RDFS reasoning is in *P*, if *G* does not contain blank nodes (ter Horst, 2002)
- RDFS reasoning is *P*-complete (i.e. in *P* but not in *NC; Patel-Schneider, 2012*)
- NC problem can be solved in O(log<sup>c</sup>n) time using O(n<sup>k</sup>) processors
  - i.e. you can trade parallelism for complexity