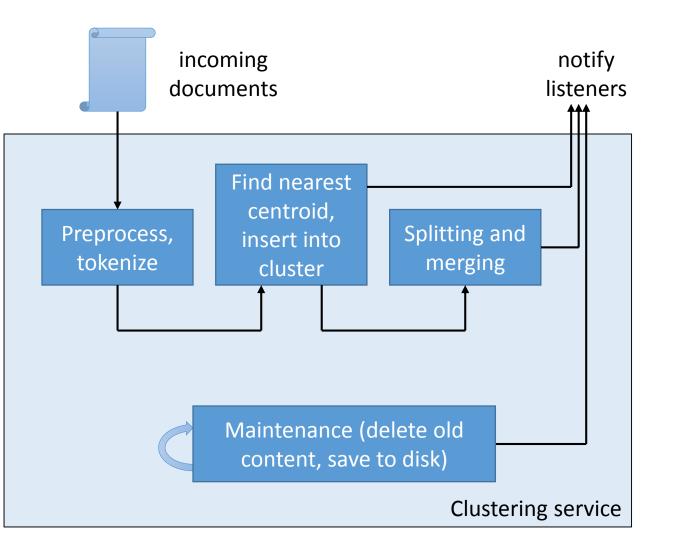
A high-performance multithreaded approach for clustering a stream of documents Janez Brank, Gregor Leban, Marko Grobelnik

#### Motivation

- Clustering
- Stream of documents
  - On-line approach
  - Asynchronous
  - Adaptive to changes in the stream
    - New clusters, splits, merges, discard old documents etc.
- Need for paralellization
  - Current aim: make good use of the parallel processing abilities of an individual computer
  - Multithreading, not distributed computing
- Main application: EventRegistry
  - Events  $\approx$  Clusters of news articles

#### Architecture

- NewsCluster web service
- Receives new documents as HTTP requests
- Notifies listeners (by making HTTP requests) about changes
  - Cluster membership
  - Cluster splits/merges
  - Cluster medians
- Background tasks
  - Delete old clusters and their documents
  - Periodically save state to disk



- Main idea:
  - Assign each document to the cluster with the nearest centroid
- We can't do too much reassigning of documents between clusters
  - Because it's an on-line setting and we don't have time
  - Because the application wants a certain amount of stability in the clusters
- Cluster maintenance consists of
  - Occasional splits/merges
  - Deleting old clusters

- In our input stream, there are many duplicates or near-duplicates
  - Basically the same article coming through several sources
  - Two documents are duplicates if they
    - Have the same title (modulo whitespace) and
    - Very similar TF-vectors (L1 distance below a threshold)
  - If a new article is found to be a duplicate of an existing article, it is discarded without further processing

- Document representation
  - Bag of words (TF-IDF vector)
  - Bag of concepts [if provided by the user]
  - Relative weight of each part in the resulting feature vector is customizable
  - Cosine similarity is used to compare feature vectors

Prepare feature vector

- The document will be assigned to the cluster whose centroid is the closest to the document
  - Compute cosine similarity between the document and the centroids of all clusters
  - If even the closest centroid is far enough, start a new cluster containing just this document

Discard obvious duplicates

Prepare feature vector

Find nearest centroid

- Each cluster maintains various statistics that are updated incrementally when the cluster changes
  - Sum of feature vectors
  - Per-feature variances
  - Medoid
- The system supports weighting documents with exponentially time-decaying weights
  - But this is currently not used in our application



Prepare feature vector

Find nearest centroid

Insert document into cluster

- After every few additions, we consider splitting the cluster into two subclusters
  - The method is based on bisecting k-means
  - Project all points onto a line and split them based on whether they fall left or right of the projection of the centroid
  - In the first pass, the line is simply the principal component, later it's a line through the centroids from the previous pass
- Accept the split if:
  - The Bayesian information criterion is met, and
  - The resulting subclusters would not meet our merge criteria
- Split by timestamp:
  - If the variance of timestamps is above a threshold
  - Split a cluster into an "older" and a "newer" subcluster
  - BIC is used to choose the relative size of the subclusters
  - Hopefully the older subcluster will be discarded soon

Discard obvious duplicates

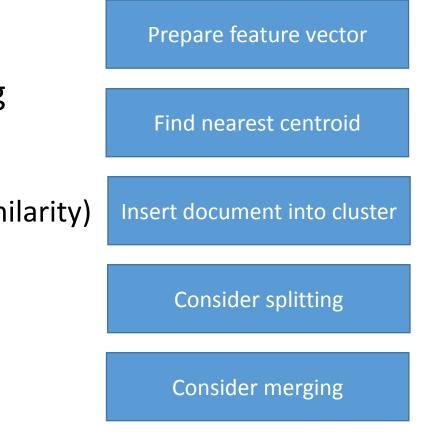
Prepare feature vector

Find nearest centroid

Insert document into cluster

Consider splitting

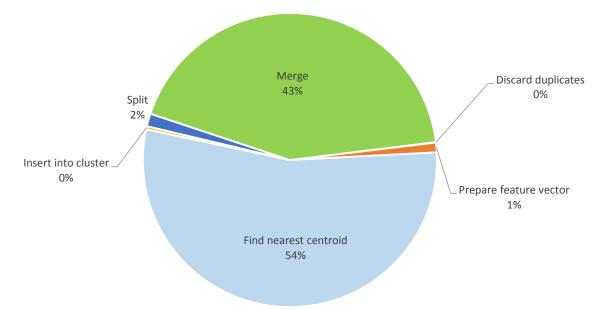
- After every few additions, we consider merging the cluster with some other cluster
  - Candidates for merging are those clusters whose centroid is the closest to our (in terms of cosine similarity)
  - Centroids tend to be dense vectors, so cosine computations are slow
    - Use "pruned" centroids for an intial filtering step
  - Accept the merge if:
    - Cosine similarity is above a threshold; or
    - Lughofer's ellipsoid overlap criterion



Discard obvious duplicates

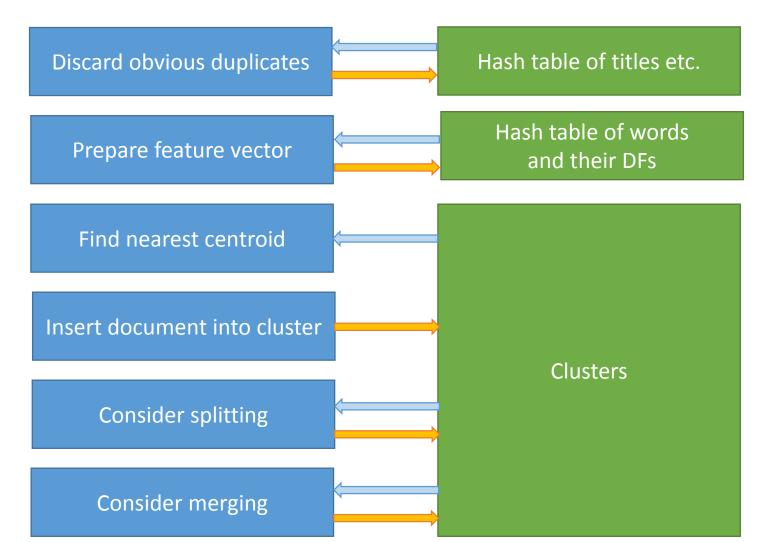
## Considerations for paralellization

- Which are the most time-consuming parts of the process?
- 54%: computing cosines between the new document and all centroids
- 43%: computing cosines between one centroid and all other centroids to find merge candidates
- 3%: everything else



## Considerations for paralellization

- Shared data structures:
  - Title hash table for duplicate detection
  - Word and DF table for bag-of-words vectors
  - Cluster membership and statistics



## Downsides of fine-grained paralellization

- Naive idea: assign each document to one worker thread
  - This thread executes the entire clustering algorithm for this document
- Downside #1: too much locking
- Downside #2: information relevant to the same cluster is scattered across multiple worker threads

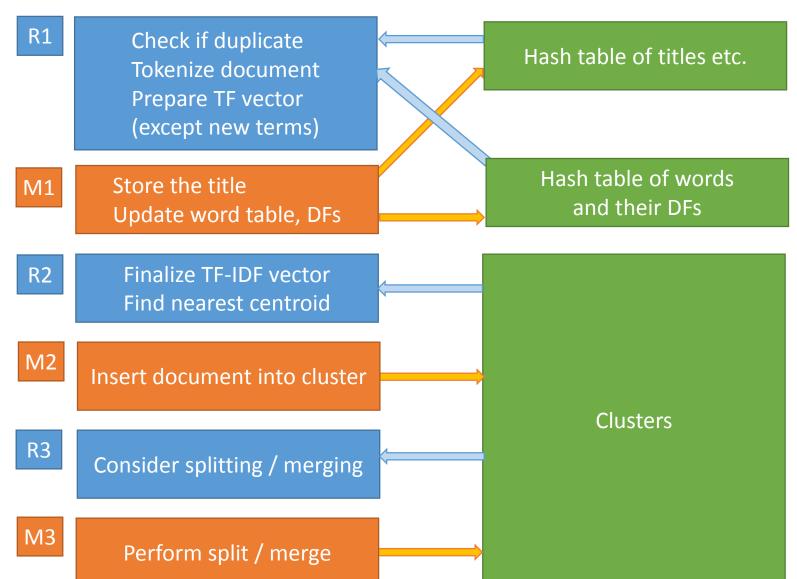
Thread 1	Thread 2	Thread 3	Thread 4
Cluster C			
Computes cosine between C and document d <sub>1</sub> (for insertion) or between C and centroid C' (for merging)	Inserts d <sub>2</sub> into C, updates centroid etc.	Trying to split <i>C</i> into <i>C</i> <sub>1</sub> and <i>C</i> <sub>2</sub>	Trying to merge <i>C</i> with <i>C</i> <sub>3</sub>

#### Avoiding cluster-level locking

- If we want to avoid per-cluster locking:
  - No thread may modify a cluster
  - While some other thread is looping through clusters (to find the nearest centroid or to find merge candidates)
    - But each thread spends 95% its time doing this
  - And yet each thread will need to modify a cluster at some point (to insert the new document into it)
    - If it has to wait for all other threads to finish looping through clusters, that's a lot of waiting
    - And it doesn't help much with the problem of having multiple threads trying to do different things to the same cluster

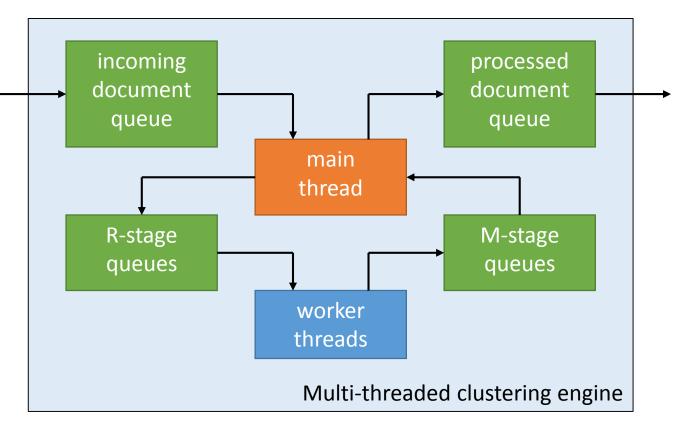
#### Read and write stages

 We can rephrase our approach a little to make it clearer which steps need to modify shared data structures



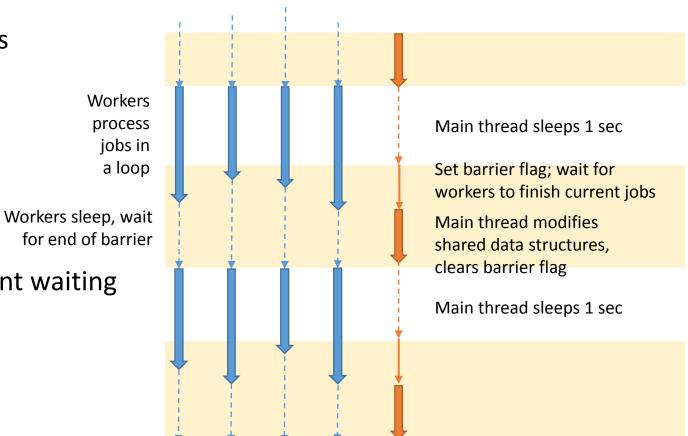
#### Main vs. Worker threads

- There is no reason why all the steps should be done by the same thread
  - *N* worker threads: perform read stages
  - 1 main thread: performs modify stages for all the requests
  - Only the main thread modifies any global data structures



#### Barrier-based parallelization

- Main thread needs to block all worker threads while it updates the shared data structures
  - Set a "barrier" flag to stop issuing new jobs to worker threads
  - Wait for all worker threads to finish their current job
  - Main thread can now modify shared data
  - Clear the barrier flag so the worker threads can resume
- We have a barrier once per sec
  - To reduce the amount of time spent waiting
  - This means each requests needs 3 seconds to be fully processed



#### Barrier processing in the main thread

- For each request that went through an R-stage since the previous barrier, we now have to perform the corresponding M-stage
  - M1: add new titles, terms to shared hash tables, update DFs
  - M3:
    - In the R3 stages, the worker threads determined whether splits/merges should be done, and how exactly the documents should be split/merged
    - The main thread now carries out these split/merge proposals, ignoring those that would clash with already-processed proposals
  - M2:
    - In the R2 stages, the worker threads recommended where to insert which new document
    - The main thread now carries out these insertions and updates cluster statistics
    - While taking into account the splits/merges just carried out in M3

## Conclusions and future work

- Multi-threaded clustering approach
  - Makes good use of parallel processing within a single computer
  - Low amount of locking and waiting
  - High throughput at the cost of high latency
  - Further paralellism comes from processing documents for different languages separately
- Possible future extensions
  - Use random projections to speed up the computation of cosine similarities
  - Hierarchical clustering
  - Distributed processing

